Mauricio Romero

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Consider the following game among 100 people. Each individual selects a number, s_i , between 20 and 60.

Let a_{-i} be the average of the number selected by the other 99 people. i.e. $a_{-i} = \sum_{j \text{ neq}i} \frac{s_j}{99}$.

The utility function of the individual i is $u_i(s_i, s_{-i}) = 100 - (s_i - \frac{3}{2}a_{-i})^2$

► Each individual maximizes his utility, FOC:

$$-2(s_i - \frac{3}{2}a_{-i}) = 0$$

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- ► Individuals would prefer to select a number that is exactly equal to 1.5 times the average of the others
- ► That is they would like to choose $s_i = \frac{3}{2}a_{-i}$
- ▶ but $a_{-i} \in [20, 60]$
- ▶ Therefore $s_i = 20$ is dominated by $s_i = 30$

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- Nowing this, all individuals believe that everyone else will select a number between 30 and 60 (i.e., $a_{-i} \in [30, 60]$)
- ▶ Playing a number between 30 and 45 (not including) would be strictly dominated by playing 45

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- Nowing this, all individuals believe that everyone else will select a number between 45 and 60 (i.e., $a_{-i} \in [45, 60]$)
- ▶ 60 would dominate any other selection and therefore all the players select 60.
- The solution by means of iterated elimination of dominated strategies is (60, 60, ..., 60)

100 times

Dominance
Weakly dominated strategies

Nash equilibrium

Some examples

Relationship to dominance

Examples
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	а	b
Α	3, 4	4, 3
В	5, 3	3, 5
С	5, 3	4, 3

► There is no strictly dominated strategy

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- ▶ It's tempting to think player 1 would never play C

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- There is no strictly dominated strategy
- ▶ However, C always gives at least the same utility to player 1 as B
- ▶ It's tempting to think player 1 would never play C
- ► However, if player 1 is sure that player two is going to play *a* he would be completely indifferent between playing *B* or *C*

Definition

 s_i weakly dominates s_i' if for all opponent pure strategy profiles, $s_{-i} \in S_{-i}$,

$$u_i(s_i,s_{-i}) \geq u_i(_i',s_{-i})$$

and there is at least one opponent strategy profile $s_{-i} \in S_{-i}$ for which

$$u_i(s_i, s_{-i}) > u_i(s_i', s_{-i}).$$

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- Given the assumptions we have, we can not eliminate a weakly dominated strategy
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► Even so, it sounds "logical" to do so and has the potential to greatly simplify a game

► There is a problem, and that is that the order in which we eliminate the strategies matters

	а	b
Α	3, 4	4, 3
В	5, 3	3, 5
С	5, 3	4, 3

▶ If we eliminate *B* (*C* dominates weakly), then *a* weakly dominates *b* and we can eliminate *b* and therefore player 1 would never play A. This leads to the result (*C*, *a*).

	а	b
Α	3, 4	4, 3
В	5, 3	3, 5
С	5, 3	4, 3

- ► If we eliminate B (C dominates weakly), then a weakly dominates b and we can eliminate b and therefore player 1 would never play A. This leads to the result (C, a).
- ▶ If on the other hand, we notice that A is also weakly dominated by C then we can eliminate it in the first round, and this would eliminate a in the second round and therefore B would be eliminated. This would result in (C, b).

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Remember the definition of competitive equilibrium in a market economy.

Definition

A competitive equilibrium in a market economy is a vector of prices and baskets x_i such that: 1) x_i maximizes the utility of each individual given the price vector i.e.

$$x_i = \arg\max_{p \ cdot x_i \le p \cdot w_i} u(x_i)$$

2) the markets empty.

$$\sum_{i} x_{i} = \sum_{i} w_{i}$$

▶ 1) means that given the prices, individuals have no incentive to demand a different amount

▶ 1) means that given the prices, individuals have no incentive to demand a different amount

▶ The idea is to extend this concept to strategic situations

Best response

We denote $BR_i(s_{-i})$ (best response) as the set of strategies of individual i that maximize her utility given that other individuals follow the strategy profile s_{-i} . Formally,

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Given a strategy profile of opponents s_{-i} , we can define the best response of player i:

$$BR_i(s_{-i}) = \underset{s_i' \in S_i}{\text{arg max}} u_i(s_i', s_{-i}).$$

▶ $s_i \in BR_i(s_{-i})$ if and only if $u_i(s_i, s_{-i}) \ge u_i(s_i', s_{-i})$ for all $s_i' \in S_i$



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$$BR_i(s_{-i}) = \underset{s_i' \in S_i}{\text{arg max}} u_i(s_i', s_{-i}).$$

- $s_i \in BR_i(s_{-i})$ if and only if $u_i(s_i, s_{-i}) \ge u_i(s_i', s_{-i})$ for all $s_i' \in S_i$
- ▶ There could be multiple strategies in $BR_i(s_{-i})$ but all such strategies give the same utility to player i if the opponents are indeed playing according to s_{-i}

Nash equilibrium

Definition

Suppose that we have a game $(I = \{1, 2, \dots, n\}, S_1, \dots, S_n, u_1, \dots, u_n)$. Then a strategy profile $s^* = (s_1^*, \dots, s_n^*)$ is a **pure strategy** Nash equilibrium if for every i and for every $s_i \in S_i$,

$$u_i(s_i^*, s_{-i}^*) \geq u_i(s_i, s_{-i}^*).$$



Nash equilibrium

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Suppose that we have a game $(I = \{1, 2, ..., n\}, S_1, ..., S_n, u_1, ..., u_n)$. Then a strategy profile $s^* = (s_1^*, ..., s_n^*)$ is a **pure strategy** Nash equilibrium if for every $i, s_i^* \in BR_i(s^*)$.

► Analogous to that of a competitive equilibrium in the sense that nobody has unilateral incentives to deviate

Nash equilibrium

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- once this equilibrium is reached, nobody has incentives to move from there

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Definition

Suppose that we have a game $(I = \{1, 2, ..., n\}, S_1, ..., S_n, u_1, ..., u_n)$. Then a strategy profile $s^* = (s_1^*, ..., s_n^*)$ is a **pure strategy** Nash equilibrium if for every $i, s_i^* \in BR_i(s_{-i}^*)$.

- Analogous to that of a competitive equilibrium in the sense that nobody has unilateral incentives to deviate
- once this equilibrium is reached, nobody has incentives to move from there
- ► This is a concept of stability, but there is no way to ensure, or predict, that the game will reach this equilibrium

Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

▶ Consider the following game among 2 people. Each individual selects a number, s_i , between 20 and 60.

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The utility function of the individual i is $u_i(s_i, s_{-i}) = 100 - (s_i - \frac{3}{2}s_{-i})^2$

The best response of an individual is given by

$$s_i(s_{-i})^* = \begin{cases} \frac{3}{2}s_{-i} & \text{si } s_{-i} \le 40\\ 60 & \text{si } s_{-i} > 40 \end{cases}$$

The Nash equilibrium is where both BR functions intersect (i.e., when both play 60)

Prisoner's dilemma

	С	NC
C	5,5	0,10
NC	10,0	2,2

Prisoner's dilemma

	С	NC
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The best response functions are:

$$BR_i(s_{-i}) = \begin{cases} NC & \text{si } s_i = C \\ NC & \text{si } s_i = NC \end{cases}$$

The Nash equilibrium is where both BR functions intersect (i.e., when both play NC, NC)

Best response of 1 to 2 playing C

	С	NC
С	5,5	0,10
NC	<u>10</u> ,0	2,2

Best response of 1 to 2 playing NC

	С	NC
С	5,5	0,10
NC	<u>10</u> ,0	<u>2</u> ,2

Best response of 2 to 1 playing C

	С	NC
С	5,5	0, <u>10</u>
NC	<u>10</u> ,0	<u>2</u> ,2

Best response of 2 to 1 playing NC

	С	NC
С	5,5	0, <u>10</u>
NC	<u>10</u> ,0	<u>2,2</u>

When underlined for both players, it is a Nash equilibrium (both are doing their BR)

Battle of the sexes

	G	Р
G	2,1	0,0
Р	0,0	1,2

Battle of the sexes

	G	Р
G	<u>2,1</u>	0,0
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$$BR_i(s_{-i}) = \begin{cases} G & \text{si } s_i = G \\ P & \text{si } s_i = P \end{cases}$$

Battle of the sexes

	G	Р
G	<u>2,1</u>	0,0
Р	0,0	<u>1,2</u>

$$BR_i(s_{-i}) = \begin{cases} G & \text{si } s_i = G \\ P & \text{si } s_i = P \end{cases}$$

Thus, (G, G) y (P, P) are both Nash equilibrium

Matching pennies (Pares o Nones I)

	1	2
1	(1000,-1000)	(-1000,1000)
2	(-1000,1000)	(1000,-1000)

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1	(<u>1000</u> ,-1000)	(-1000, <u>1000</u>)
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Matching pennies (Pares o Nones I)

	1	2
1	(<u>1000</u> ,-1000)	(-1000, <u>1000</u>)
2	(-1000, <u>1000</u>)	(<u>1000</u> ,-1000)

$$BR_1(s_2) = \begin{cases} 1 & \text{si } s_2 = 1 \\ 2 & \text{si } s_2 = 2 \end{cases}$$
 $BR_2(s_1) = \begin{cases} 2 & \text{si } s_1 = 1 \\ 1 & \text{si } s_2 = 2 \end{cases}$

There is no Nash equilibrium in pure strategies

Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Nash equilibrium survive IDSDS

Theorem

Every Nash equilibrium survives the iterative elimination of strictly dominated strategies

By contradiction:

► Suppose it is not true

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- It must have been that

$$u_i(s_i^*, s_{-i}) < u_i(s_i, s_{-i}) \forall s_{-i} \in S_{-i}$$

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► In particular

$$u_i(s_i^*, s_{-i}^*) < u_i(s_i, s_{-i}^*)$$

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▶ But this means s_i^* is not the best response of individual i to s_{-i}^*

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- ▶ But this means s_i^* is not the best response of individual i to s_{-i}^*
- And this is a contradiction!



Nash equilibrium survive IDSDS

Theorem

If the process of IDSDS comes to a single solution, that solution is a Nash Equilibrium and is unique.

First let's proof its a Nash Equilibrium. The fact that is unique is trivial by the previous theorem.

Proof.

By contradiction:

Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium

First let's proof its a Nash Equilibrium. The fact that is unique is trivial by the previous theorem.

Proof.

- Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium
- For some individual i there exits s_i such that

$$u_i(s_i, s_{-i}^*) > u_i(s_i^*, s_{-i}^*)$$

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Proof.

By contradiction:

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But then s_i could not have been eliminated



Proof

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Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Lecture 12: Game Theory // Nash equilibrium

Dominance

Nash equilibrium

Some examples

Relationship to dominance

Examples

Lecture 12: Game Theory // Nash equilibrium

Dominance

Weakly dominated strategies

Nash equilibrium

Some examples

Relationship to dominance

Examples

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$$P(Q) = 120 - Q, Q = q_1 + q_2.$$

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- Suppose that there are two firms that produce the same product have zero marginal cost of production.
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$$P(Q) = 120 - Q, Q = q_1 + q_2.$$

- ▶ Strategy space is $S_i = [0, +\infty)$
- ▶ The utility function of player *i* is given by:

$$\pi_1(q_1, q_2) = (120 - (q_1 + q_2))q_1,$$

 $\pi_2(q_1, q_2) = (120 - (q_1 + q_2))q_2.$

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$$\frac{d\pi_i}{dq_i}(120-q_i-q_{-i})q_i=120-2q_i-q_{-i}$$

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▶ Therefore 60 strictly dominates any $q_i \in (60, 120]$

$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}.$$

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▶ for any $q_i \in [0,60]$, there exists some $q_{-i} \in [0,+\infty)$ such that $BR_i(q_{-i}) = q_i$

$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}.$$

- ▶ for any $q_i \in [0, 60]$, there exists some $q_{-i} \in [0, +\infty)$ such that $BR_i(q_{-i}) = q_i$
- Such a q_i can never be strictly dominated

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- ▶ for any $q_i \in [0,60]$, there exists some $q_{-i} \in [0,+\infty)$ such that $BR_i(q_{-i}) = q_i$
- Such a q_i can never be strictly dominated
- After one round of deletion of strictly dominated strategies, we are left with: $S_i = [0, 60]$

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- $ightharpoonup q_{-i} = [0,60]$
- ▶ Therefore $q_i \in [0,30)$ are strictly dominated by $q_i = 30$

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- $ightharpoonup q_{-i} = [0,60]$
- ▶ Therefore $q_i \in [0,30)$ are strictly dominated by $q_i = 30$
- After two rounds of deletion of strictly dominated strategies, we are left with: $S_i = [30, 60]$

$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}.$$

- $ightharpoonup q_{-i} = [30, 60]$
- ▶ 45 strictly dominates all strategies $q_i \in (45, 60]$
- After three rounds of deletion of strictly dominated strategies, we are left with: $S_i = [30, 45]$

$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}.$$

- $ightharpoonup q_{-i} = [30, 45]$
- ▶ 37.5 strictly dominates all strategies $q_i \in [37.5, 45]$
- After four rounds of deletion of strictly dominated strategies, we are left with: $S_i = [30, 45]$

After (infinitely) many iterations, the only remaining strategies are $S_i = 40$

▶ The unique solution by IDSDS is $q_1^* = q_2^* = 40$.

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At any Nash equilibrium, we must have: $q_1^* \in BR_1(q_2^*)$ and $q_2^* \in BR_2(q_1^*)$.

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$$q_1^* = \frac{120 - q_2^*}{2}, q_2^* = \frac{120 - q_1^*}{2}.$$

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At any Nash equilibrium, we must have: $q_1^* \in BR_1(q_2^*)$ and $q_2^* \in BR_2(q_1^*)$.

$$q_1^* = \frac{120 - q_2^*}{2}, q_2^* = \frac{120 - q_1^*}{2}.$$

▶ We can solve for q_1^* and q_2^* to obtain:

$$q_1^* = 40, q_2^* = 40, Q^* = 80, \Pi_1^* = \Pi_2^* = 1600.$$

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- ▶ In a duopoly, externalities are imposed on the other firm

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$$\max_{q_i} p(q_i + Q_{-i})q_i - c_i(q_i)$$



First order condition implies:

$$q_i rac{dP}{dQ}(q_i + Q_{-i}) + P(q_i + Q_{-i}) = rac{dc_i}{dq_i}(q_i)$$
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Therefore in a pure strategy Nash equilibrium $(q_1^*, q_2^*, \dots, q_n^*)$ with $Q^* = q_1^* + q_2^* + \dots + q_n^*$, we must have:

$$\begin{split} \frac{P(Q^*) - \frac{dc_1}{dq_1}(q_1^*)}{P(Q^*)} &= -\frac{q_1^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}, \\ \frac{P(Q^*) - \frac{dc_2}{dq_2}(q_2^*)}{P(Q^*)} &= -\frac{q_2^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}, \\ &\vdots \\ \frac{P(Q^*) - \frac{dc_n}{dq_n}(q_n^*)}{P(Q^*)} &= -\frac{q_n^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}. \end{split}$$

▶ Suppose that all firms have exactly the same cost function *c*

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Rewriting

$$P(Q^*) = \frac{1}{1 + \frac{1}{n} \frac{1}{\varepsilon_{Q,P}(Q^*)}} \frac{\partial c}{\partial q} \left(\frac{Q^*}{n} \right).$$



Lecture 12: Game Theory // Nash equilibrium

Dominance

Weakly dominated strategies

Nash equilibrium

Some examples

Relationship to dominance

Examples

Cournot Competition

Cartels

- ▶ Suppose there are three firms who face zero marginal cost
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In a Nash equilibrium we must have:

$$q_1^* = \frac{1 - q_2^* - q_3^*}{2}$$

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- ▶ Price is $p^* = 1/4$ and all firms get the same profits of 1/16



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▶ Therefore

$$q_A^* = rac{1-q_B^*}{2} \ q_B^* = rac{1-q_A^*}{2}.$$

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- ► Firm 3 is obtaining a disproportionate share of the joint profits (more than 1/3)

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- Firm 3 clearly wants to stay out

► This exercise explains the difficulties associated with sustaining collusive agreements (e.g., the OPEC cartel)