Mauricio Romero

Mixed strategies

Examples

Mixed strategies

Examples

# Consider rock/paper/scissors

|          | Rock | Paper | Scissors |
|----------|------|-------|----------|
| Rock     | 0,0  | -1,1  | 1,-1     |
| Paper    | 1,-1 | 0,0   | -1,1     |
| Scissors | -1,1 | 1,-1  | 0,0      |

► This game is entirely stochastic (ability has nothing to do with your chances of winning)

## Consider rock/paper/scissors

|          | Rock | Paper | Scissors |
|----------|------|-------|----------|
| Rock     | 0,0  | -1,1  | 1,-1     |
| Paper    | 1,-1 | 0,0   | -1,1     |
| Scissors | -1,1 | 1,-1  | 0,0      |

- ► This game is entirely stochastic (ability has nothing to do with your chances of winning)
- ▶ The probability of winning with every strategy is the same

# Consider rock/paper/scissors

|          | Rock | Paper | Scissors |
|----------|------|-------|----------|
| Rock     | 0,0  | -1,1  | 1,-1     |
| Paper    | 1,-1 | 0,0   | -1,1     |
| Scissors | -1,1 | 1,-1  | 0,0      |

- ► This game is entirely stochastic (ability has nothing to do with your chances of winning)
- The probability of winning with every strategy is the same
- Thus, people tend choose randomly which of the three options to play

## Consider rock/paper/scissors

|          | Rock | Paper | Scissors |
|----------|------|-------|----------|
| Rock     | 0,0  | -1,1  | 1,-1     |
| Paper    | 1,-1 | 0,0   | -1,1     |
| Scissors | -1,1 | 1,-1  | 0,0      |

- ► This game is entirely stochastic (ability has nothing to do with your chances of winning)
- The probability of winning with every strategy is the same
- ▶ Thus, people *tend* choose randomly which of the three options to play
- ▶ We would like the concept of Nash equilibrium to reflect this

### Definition

A mixed strategy  $\sigma_i$  is a function  $\sigma_i: S_i \to [0,1]$  such that

$$\sum_{s_i \in S_i} \sigma_i(s_i) = 1.$$

 $ightharpoonup \sigma_i(s_i)$  represents the probability with which player i plays  $s_i$ 

### Definition

A mixed strategy  $\sigma_i$  is a function  $\sigma_i: S_i \to [0,1]$  such that

$$\sum_{s_i \in S_i} \sigma_i(s_i) = 1.$$

- $ightharpoonup \sigma_i(s_i)$  represents the probability with which player i plays  $s_i$
- ▶ A **pure strategy** is simply a mixed strategy  $\sigma_i$  that plays some strategy  $s_i \in S_i$  with probability one

### Definition

A mixed strategy  $\sigma_i$  is a function  $\sigma_i:S_i\to [0,1]$  such that

$$\sum_{s_i \in S_i} \sigma_i(s_i) = 1.$$

- $ightharpoonup \sigma_i(s_i)$  represents the probability with which player i plays  $s_i$
- ▶ A **pure strategy** is simply a mixed strategy  $\sigma_i$  that plays some strategy  $s_i \in S_i$  with probability one
- $\blacktriangleright$  We will denote the set of all mixed strategies of player i by  $\Sigma_i$

▶ Given a mixed strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$ , we need a way to define how players evaluate payoffs of mixed strategy profiles

▶ Given a mixed strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$ , we need a way to define how players evaluate payoffs of mixed strategy profiles

$$u_1(\sigma_1,\sigma_2,\ldots,\sigma_n)=\sum_{s\in S}u_1(s_1,s_2,\ldots,s_n)\sigma_1(s_1)\sigma_2(s_2)\cdots\sigma_n(s_n).$$

▶ Given a mixed strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$ , we need a way to define how players evaluate payoffs of mixed strategy profiles

$$u_1(\sigma_1,\sigma_2,\ldots,\sigma_n)=\sum_{s\in S}u_1(s_1,s_2,\ldots,s_n)\sigma_1(s_1)\sigma_2(s_2)\cdots\sigma_n(s_n).$$

For instance, assume my opponent is playing randomizing over paper and scissors with probability  $\frac{1}{2}$  (i.e.,  $\sigma_{-i} = (0, \frac{1}{2}, \frac{1}{2})$ )

▶ Given a mixed strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$ , we need a way to define how players evaluate payoffs of mixed strategy profiles

$$u_1(\sigma_1,\sigma_2,\ldots,\sigma_n)=\sum_{s\in S}u_1(s_1,s_2,\ldots,s_n)\sigma_1(s_1)\sigma_2(s_2)\cdots\sigma_n(s_n).$$

- For instance, assume my opponent is playing randomizing over paper and scissors with probability  $\frac{1}{2}$  (i.e.,  $\sigma_{-i} = (0, \frac{1}{2}, \frac{1}{2})$ )
- ▶ The expected utility of playing "rock" is

$$E(U_i(rock, \sigma_{-i})) = -1\frac{1}{2} + 1\frac{1}{2} = 0$$

▶ Given a mixed strategy profile  $(\sigma_1, \sigma_2, \dots, \sigma_n)$ , we need a way to define how players evaluate payoffs of mixed strategy profiles

$$u_1(\sigma_1,\sigma_2,\ldots,\sigma_n)=\sum_{s\in S}u_1(s_1,s_2,\ldots,s_n)\sigma_1(s_1)\sigma_2(s_2)\cdots\sigma_n(s_n).$$

- For instance, assume my opponent is playing randomizing over paper and scissors with probability  $\frac{1}{2}$  (i.e.,  $\sigma_{-i} = (0, \frac{1}{2}, \frac{1}{2})$ )
- The expected utility of playing "rock" is

$$E(U_i(rock, \sigma_{-i})) = -1\frac{1}{2} + 1\frac{1}{2} = 0$$

If I'm randomizing over rock and scissors (i.e.,  $\sigma_i = (\frac{1}{2}, 0, \frac{1}{2})$ ) then

$$E(U_i(\sigma,\sigma_{-i})) = \underbrace{-1\frac{1}{4}}_{} + \underbrace{1\frac{1}{4}}_{} + \underbrace{1\frac{1}{4}}_{} + \underbrace{0\frac{1}{4}}_{} = \frac{1}{4}$$

### Definition

A (possibly mixed) strategy profile  $(\sigma_1^*, \sigma_2^*, \dots, \sigma_n)^*$  is a Nash equilibrium if and only if for every i,

$$u_i(\sigma_i^*, \sigma_{-i}^*) \geq u_i(\sigma_i, \sigma_{-i}^*)$$

for all  $\sigma_i \in \Sigma_i$ .

# Definition (Mixed Strategy Dominance Definition A)

Let  $\sigma_i, \sigma_i'$  be two mixed strategies of player i. Then  $\sigma_i$  strictly dominates  $\sigma_i'$  if for all mixed strategies of the opponents,  $\sigma_{-i}$ ,

$$u_i(\sigma_i, \sigma_{-i}) > u_i(\sigma'_i, \sigma_{-i}).$$

If  $\sigma_i$  is better than  $\sigma_i'$  no matter what **pure strategy** opponents play, then  $\sigma_i$  is also strictly better than  $\sigma_i'$  no matter what **mixed strategies** opponents play

### **Theorem**

Let  $\sigma_i$  and  $\sigma_i'$  be two mixed strategies of player i. Then  $\sigma_i$  strictly dominates  $\sigma_i'$  if and only if for all  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

## Proof- Part 1

▶ Since  $S_{-i} \subseteq \Sigma_{-i}$ , if  $\sigma_i$  strictly dominates  $\sigma'_i$ 

### Proof- Part 1

▶ Since 
$$S_{-i} \subseteq \Sigma_{-i}$$
, if  $\sigma_i$  strictly dominates  $\sigma'_i$ 

▶ Then for all  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

### Proof - Part 2

▶ To prove the other direction, suppose that for all  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

### Proof - Part 2

▶ To prove the other direction, suppose that for all  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

▶ For any  $\sigma_{-i}$ ,

$$u_i(\sigma_i, \sigma_{-i}) = \sum_{s_i \in S_i} \sum_{s_{-i} \in S_{-i}} \sigma_i(s_i) \sigma_{-i}(s_{-i}) u_i(s_i, s_{-i})$$

$$= \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i}) \sum_{s_i \in S_i} \sigma_i(s_i) u_i(s_i, s_{-i})$$

$$= \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i}) u_i(\sigma_i, s_{-i})$$

### Proof - Part 2

▶ To prove the other direction, suppose that for all  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

▶ For any  $\sigma_{-i}$ ,

$$u_i(\sigma_i, \sigma_{-i}) = \sum_{s_i \in S_i} \sum_{s_{-i} \in S_{-i}} \sigma_i(s_i) \sigma_{-i}(s_{-i}) u_i(s_i, s_{-i})$$

$$= \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i}) \sum_{s_i \in S_i} \sigma_i(s_i) u_i(s_i, s_{-i})$$

$$= \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i}) u_i(\sigma_i, s_{-i})$$

► So

$$u_i(\sigma_i, \sigma_{-i}) = \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i})u_i(\sigma_i, s_{-i}) > \sum_{s_{-i} \in S_{-i}} \sigma_{-i}(s_{-i})u_i(\sigma_i', s_{-i}) = u_i(\sigma_i', \sigma_{-i})$$

# Definition (Mixed Strategy Dominance Definition B)

Let  $\sigma_i, \sigma_i'$  be two mixed strategies of player i. Then  $\sigma_i$  strictly dominates  $\sigma_i'$  if for all pure strategies of the opponents,  $s_{-i} \in S_{-i}$ ,

$$u_i(\sigma_i, s_{-i}) > u_i(\sigma'_i, s_{-i}).$$

Mixed strategies

Examples

Mixed strategies

Examples

|   | G   | Р   |
|---|-----|-----|
| G | 2,1 | 0,0 |
| Р | 0,0 | 1,2 |

|   | G          | Р          |
|---|------------|------------|
| G | <u>2,1</u> | 0,0        |
| Р | 0,0        | <u>1,2</u> |

▶ There are two pure strategy equilibria (G, G) and (P, P)

|   | G          | Р          |
|---|------------|------------|
| G | <u>2,1</u> | 0,0        |
| Р | 0,0        | <u>1,2</u> |

▶ There are two pure strategy equilibria (G, G) and (P, P)

We now look for Nash equilibria that involve randomization by the players

Let  $\lambda$  be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

Let  $\lambda$  be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

$$u_1(\lambda,q)=2\lambda q+(1-\lambda)(1-q).$$

- Let  $\lambda$  be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

$$u_1(\lambda,q)=2\lambda q+(1-\lambda)(1-q).$$

▶ Case 1: If q>1/3, then 2q>2/3>1-q and therefore, the best response is  $\lambda=1$ 

- Let  $\lambda$  be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

$$u_1(\lambda,q)=2\lambda q+(1-\lambda)(1-q).$$

- ▶ Case 1: If q > 1/3, then 2q > 2/3 > 1 q and therefore, the best response is  $\lambda = 1$
- ▶ Case 2: if q = 1/3, then 2q = 2/3 = 1 q and therefore, the best response is  $\lambda \in [0, 1]$

- Let λ be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

$$u_1(\lambda,q)=2\lambda q+(1-\lambda)(1-q).$$

- ▶ Case 1: If q > 1/3, then 2q > 2/3 > 1 q and therefore, the best response is  $\lambda = 1$
- ▶ Case 2: if q = 1/3, then 2q = 2/3 = 1 q and therefore, the best response is  $\lambda \in [0, 1]$
- ▶ Case 3: If q < 1/3, then 2q < 2/3 < 1-q and therefore the best response is  $\lambda = 0$

Let λ be the probability with which player 1 chooses G and q be the probability with which player 2 plays G

$$u_1(\lambda, q) = 2\lambda q + (1 - \lambda)(1 - q).$$

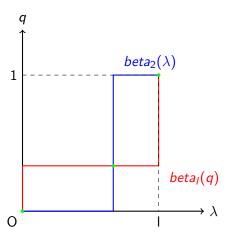
- ▶ Case 1: If q > 1/3, then 2q > 2/3 > 1 q and therefore, the best response is  $\lambda = 1$
- ▶ Case 2: if q = 1/3, then 2q = 2/3 = 1 q and therefore, the best response is  $\lambda \in [0, 1]$
- ▶ Case 3: If q < 1/3, then 2q < 2/3 < 1 q and therefore the best response is  $\lambda = 0$
- Thus, the best response function is given by:

$$BR_1(q) = egin{cases} 1 & \text{if } q > 1/3 \\ [0,1] & \text{if } q = 1/3 \\ 0 & \text{if } q < 1/3. \end{cases}$$

Similarly we can calculate the best response function for player 2 and we get:

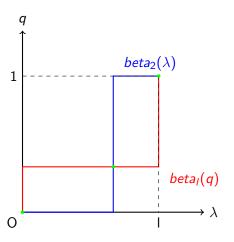
$$BR_2(\lambda) = \begin{cases} 1 & \text{if } \lambda > 2/3\\ [0,1] & \text{if } \lambda = 2/3\\ 0 & \text{if } \lambda < 2/3. \end{cases}$$

### Battle of the sexes



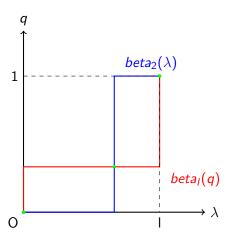
There are three points where the best response curves cross:  $(1,1),(0,0,),(\frac{2}{3},\frac{1}{3})$ 

#### Battle of the sexes



- There are three points where the best response curves cross:  $(1,1),(0,0,),(\frac{2}{3},\frac{1}{3})$
- First two are the pure strategy NE we had found before

### Battle of the sexes



- There are three points where the best response curves cross:  $(1,1),(0,0,),(\frac{2}{3},\frac{1}{3})$
- First two are the pure strategy NE we had found before
- Last is a strictly mixed NE: both players randomize



# Consider the following game

|   | Е     | F    | G    |
|---|-------|------|------|
| Α | 5, 10 | 5, 3 | 3, 4 |
| В | 1, 4  | 7, 2 | 7, 6 |
| С | 4, 2  | 8, 4 | 3, 8 |
| D | 2, 4  | 1, 3 | 8, 4 |

► Consider  $\sigma_1 = (\frac{1}{3}, \frac{1}{4}, \frac{1}{4}, \frac{1}{6})$ 

► Consider 
$$\sigma_1 = (\frac{1}{3}, \frac{1}{4}, \frac{1}{4}, \frac{1}{6})$$

$$\blacksquare U(E, \sigma_1) = 10\frac{1}{3} + 4\frac{1}{4} + 2\frac{1}{4} + 4\frac{1}{6} = 5.5$$

► Consider 
$$\sigma_1 = (\frac{1}{3}, \frac{1}{4}, \frac{1}{4}, \frac{1}{6})$$

$$\mathbb{E}U(E,\sigma_1) = 10\frac{1}{3} + 4\frac{1}{4} + 2\frac{1}{4} + 4\frac{1}{6} = 5.5$$

$$\mathbb{E}U(F,\sigma_1) = 3\frac{1}{3} + 2\frac{1}{4} + 4\frac{1}{4} + 3\frac{1}{6} = 3$$

► Consider 
$$\sigma_1 = (\frac{1}{3}, \frac{1}{4}, \frac{1}{4}, \frac{1}{6})$$

$$\mathbb{E}U(E,\sigma_1) = 10\frac{1}{3} + 4\frac{1}{4} + 2\frac{1}{4} + 4\frac{1}{6} = 5.5$$

$$\mathbb{E}U(F,\sigma_1) = 3\frac{1}{3} + 2\frac{1}{4} + 4\frac{1}{4} + 3\frac{1}{6} = 3$$

$$\mathbb{E}U(G,\sigma_1) = 4\frac{1}{3} + 6\frac{1}{4} + 8\frac{1}{4} + 4\frac{1}{6} = 5.5$$

- ► Consider  $\sigma_1 = (\frac{1}{3}, \frac{1}{4}, \frac{1}{4}, \frac{1}{6})$
- $\mathbb{E}U(E,\sigma_1) = 10\frac{1}{3} + 4\frac{1}{4} + 2\frac{1}{4} + 4\frac{1}{6} = 5.5$
- $\mathbb{E}U(F,\sigma_1) = 3\frac{1}{3} + 2\frac{1}{4} + 4\frac{1}{4} + 3\frac{1}{6} = 3$
- $\mathbb{E}U(G,\sigma_1) = 4\frac{1}{3} + 6\frac{1}{4} + 8\frac{1}{4} + 4\frac{1}{6} = 5.5$
- ▶ Then  $BR_2(\sigma_1) = \{(p, 0, 1 p), p \in [0, 1]\}$

► *G* dominates *F* (player 2)

► G dominates F (player 2)

► D dominates B (player 1)

# Reduced game

|   | Е     | G    |
|---|-------|------|
| Α | 5, 10 | 3, 4 |
| С | 4, 2  | 3, 8 |
| D | 2, 4  | 8, 4 |

Note that 
$$\sigma_1 = (p, 0, 1 - p)$$
 with  $p > \frac{2}{3}$  dominates  $C$ 

$$ightharpoonup \mathbb{E} U(\sigma_1, E) = 5p + 2(1-p) = 3p + 2$$

$$\triangleright$$
  $\mathbb{E}U(\sigma_1, G) = 3p + 8(1-p) = 8 - 5p$ 

$$\mathbb{E}U(\sigma_1, E) > U(C, E)$$

$$3p + 2 > 4$$

$$p > \frac{2}{3}$$

$$\mathbb{E}U(\sigma_1, G) > \mathbb{E}U(C, G)$$

$$8 - 3p > 3$$

$$p < \frac{5}{5} = 1$$

### Reduced game

|   | Е     | G    |
|---|-------|------|
| Α | 5, 10 | 3, 4 |
| D | 2, 4  | 8, 4 |

Lets find  $BR_1(\sigma_2 = (q, 1-q))$ 

- Lets find  $BR_1(\sigma_2 = (q, 1-q))$
- $\triangleright$   $\mathbb{E}U(A, \sigma_2) = 5q + 3(1-q) = 2q + 3$

- ▶ Lets find  $BR_1(\sigma_2 = (q, 1-q))$
- $\triangleright$   $\mathbb{E}U(A, \sigma_2) = 5q + 3(1-q) = 2q + 3$
- $ightharpoonup \mathbb{E}U(D, \sigma_2) = 2q + 8(1-q) = 8 6q$

- Lets find  $BR_1(\sigma_2 = (q, 1-q))$
- $\triangleright$   $\mathbb{E}U(A, \sigma_2) = 5q + 3(1-q) = 2q + 3$
- $ightharpoonup \mathbb{E}U(D, \sigma_2) = 2q + 8(1-q) = 8 6q$
- ▶ 8 6q > 2q + 3 if  $\frac{5}{8} > q$

- Lets find  $BR_1(\sigma_2 = (q, 1-q))$
- $\triangleright$   $\mathbb{E}U(A, \sigma_2) = 5q + 3(1-q) = 2q + 3$
- $ightharpoonup \mathbb{E}U(D, \sigma_2) = 2q + 8(1-q) = 8 6q$
- ▶ 8 6q > 2q + 3 if  $\frac{5}{8} > q$
- ▶ 8 6q < 2q + 3 if  $\frac{5}{8} < q$

Lets find 
$$BR_1(\sigma_2 = (q, 1-q))$$

$$ightharpoonup \mathbb{E}U(A, \sigma_2) = 5q + 3(1-q) = 2q + 3$$

$$ightharpoonup \mathbb{E}U(D, \sigma_2) = 2q + 8(1-q) = 8 - 6q$$

▶ 
$$8 - 6q > 2q + 3$$
 if  $\frac{5}{8} > q$ 

▶ 
$$8 - 6q < 2q + 3$$
 if  $\frac{5}{8} < q$ 

► Thus

$$BR_1(q, 1-q) = egin{cases} \sigma_1 = (0,1) & ext{if } 0 \leq q < rac{5}{8} \ \sigma_1 = (1,0) & ext{if } rac{5}{8} < q \leq 1 \ \sigma_1 = (p,1-p) & ext{if } rac{5}{8} = q \end{cases}$$

▶ Lets find  $BR_2(\sigma_1 = (p, 1-p))$ 

- Lets find  $BR_2(\sigma_1 = (p, 1-p))$
- $\blacktriangleright$   $\mathbb{E}U(\sigma_1, E) = 10p + 4(1-p) = 6p + 4$

- ▶ Lets find  $BR_2(\sigma_1 = (p, 1-p))$
- $\triangleright$   $\mathbb{E}U(\sigma_1, E) = 10p + 4(1-p) = 6p + 4$
- $ightharpoonup \mathbb{E} U(\sigma_1, G) = 4p + 4(1-p) = 4$

- ▶ Lets find  $BR_2(\sigma_1 = (p, 1-p))$
- $\triangleright$   $\mathbb{E}U(\sigma_1, E) = 10p + 4(1-p) = 6p + 4$

$$ightharpoonup \mathbb{E} U(\sigma_1, G) = 4p + 4(1-p) = 4$$

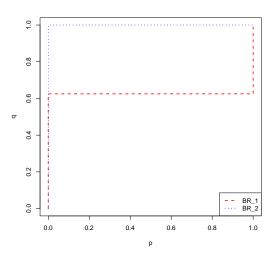
▶ 
$$6p + 4 > 4$$
 if  $p > 0$ 

- ▶ Lets find  $BR_2(\sigma_1 = (p, 1-p))$
- $\triangleright$   $\mathbb{E}U(\sigma_1, E) = 10p + 4(1-p) = 6p + 4$
- $ightharpoonup \mathbb{E} U(\sigma_1, G) = 4p + 4(1-p) = 4$
- ▶ 6p + 4 > 4 if p > 0
- ▶ 6p + 4 < 4 if p < 0.

- Lets find  $BR_2(\sigma_1 = (p, 1-p))$
- $\blacksquare U(\sigma_1, E) = 10p + 4(1-p) = 6p + 4$
- $\triangleright$   $\mathbb{E}U(\sigma_1,G) = 4p + 4(1-p) = 4$
- ▶ 6p + 4 > 4 if p > 0
- ▶ 6p + 4 < 4 if p < 0.
- ► Thus

$$BR_2(p, 1-p) = \begin{cases} \sigma_2 = (1, 0) & \text{if } p > 0 \\ \sigma_2 = (q, 1-q) & \text{if } p = 0 \end{cases}$$

### Best responses



$$\mathit{NE} = \{(D,G),(A,\sigma_2^q)\}$$
 where  $\sigma_2^q = (q,1-q)$  and  $0 \leq q \leq \frac{5}{8}$