



Lecture 12: Game Theory // Nash equilibrium

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Lecture 12: Game Theory // Nash equilibrium

- Dominance
- Nash equilibrium
- Some examples
- Relationship to dominance
- Examples

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Beauty contest

- Consider the following game among 100 people. Each individual selects a number, s_i , between 20 and 60.

- Let \bar{s} be the average of the number selected by the other 99 people. i.e.

$$\bar{s} = \frac{1}{99} \sum_{j \neq i} s_j$$

- The utility function of the individual j is $u_j(s_j, s_{-j}) = 100 - (s_j - \bar{s}_{-j})^2$

observo:
Asume s_k de los demás

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Beauty contest

- Each individual maximizes his utility. FOC:

$$-2(s_j - \bar{s}_{-j}) = 0$$

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Beauty contest

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$$-2(s - \frac{3}{2}x_i) = 0$$
- Individuals would prefer to select a number that is exactly equal to 1.5 times the average of the others

Beauty contest

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- Individuals would prefer to select a number that is exactly equal to 1.5 times the average of the others
- That is they would like to choose $s_i = \frac{3}{2}x_i$
- but $x_i \in [20, 60]$ $\frac{3}{2}x_i \in [\frac{3}{2} \cdot 20, \frac{3}{2} \cdot 60] = [30, 90]$

ai

Beauty contest

- Each individual maximizes his utility, FOC:

$$-2(s - \frac{3}{2}x_i) = 0$$
- Individuals would prefer to select a number that is exactly equal to 1.5 times the average of the others
- That is they would like to choose $s_i = \frac{3}{2}x_i$
- but $x_i \in [20, 60]$
- Therefore $s_i = 20$ is dominated by $x_i = 30$

Beauty contest

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Beauty contest

- The same goes for any number between 20 (inclusive) and 30 (not included)
- Knowing this, all individuals believe that everyone else will select a number between 30 and 60 (i.e., $a_i \in [30, 60]$)
- Playing a number between 30 and 45 (not including) would be strictly dominated by playing 45

$a_i \in [30, 60]$
 $\frac{3}{2} a_i \in [45, 90]$

Beauty contest

- The same goes for any number between 20 (inclusive) and 30 (not included)
- Knowing this, all individuals believe that everyone else will select a number between 30 and 60 (i.e., $a_i \in [30, 60]$)
- Playing a number between 30 and 45 (not including) would be strictly dominated by playing 45
- Knowing this, all individuals believe that everyone else will select a number between 45 and 60 (i.e., $a_i \in [45, 60]$)

$\frac{3}{2} a_i \in [\frac{3}{2} \cdot 45, \frac{3}{2} \cdot 60]$
 $\frac{3}{2} a_i \in [67.5, 90]$

Beauty contest

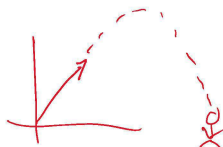
- The same goes for any number between 20 (inclusive) and 30 (not included)
- Knowing this, all individuals believe that everyone else will select a number between 30 and 60 (i.e., $a_i \in [30, 60]$)
- Playing a number between 30 and 45 (not including) would be strictly dominated by playing 45
- Knowing this, all individuals believe that everyone else will select a number between 45 and 60 (i.e., $a_i \in [45, 60]$)
- 60 would dominate any other selection and therefore all the players select 60.

Beauty contest

- The same goes for any number between 20 (inclusive) and 30 (not included)
- Knowing this, all individuals believe that everyone else will select a number between 30 and 60 (i.e., $a_i \in [30, 60]$)
- Playing a number between 30 and 45 (not including) would be strictly dominated by playing 45
- Knowing this, all individuals believe that everyone else will select a number between 45 and 60 (i.e., $a_i \in [45, 60]$)
- 60 would dominate any other selection and therefore all the players select 60.

The solution by means of iterated elimination of dominated strategies is (60, 60, ..., 60)

(20, 60)



$S = (60, 60, \dots, 60)$ no es O.P.

P.C \rightarrow MAX $100 - (s_i - \frac{3}{2} a_i)^2$

s.t. $100 - (s_2 - \frac{3}{2} a_2)^2 \geq \bar{U}_2 : \lambda_2$
 $100 - (s_3 - \frac{3}{2} a_3)^2 \geq \bar{U}_3 : \lambda_3$

\dots
 $100 - (s_{100} - \frac{3}{2} a_{100})^2 \geq \bar{U}_{100} : \lambda_{100}$

$a_i = \frac{s_1 + s_2 + \dots + s_{100}}{99}$

$\frac{\partial Y}{\partial s_1} = -2(s_1 - \frac{3}{2} a_1) - \lambda_2 (-2(s_2 - \frac{3}{2} a_2) (\frac{-3}{2} \cdot \frac{1}{99})) = 0$

$\frac{\partial Y}{\partial s_2} = -2(s_2 - \frac{3}{2} a_2) (\frac{-3}{2} \cdot \frac{1}{99}) + \lambda_3 (-2(s_3 - \frac{3}{2} a_3)) = 0$

\dots
 $\frac{\partial Y}{\partial s_{100}} = -2(s_{100} - \frac{3}{2} a_{100}) (\frac{-3}{2} \cdot \frac{1}{99}) + \lambda_{100} (-2(s_{100} - \frac{3}{2} a_{100})) = 0$

$(-2(s - \frac{3}{2} s) - \lambda_2 (-2(s - \frac{3}{2} s) (\frac{-3}{2 \cdot 99}))) = 0 \rightarrow s =$

$-2(s - \frac{3}{2} s) = \lambda_2 (-2(s - \frac{3}{2} s) (\frac{-3}{2 \cdot 99}))$

$\lambda_2 = \frac{2 \cdot 99}{-2}$

Lecture 12: Game Theory // Nash equilibrium

Dominance

Weakly dominated strategies

Nash equilibrium

Some examples

Relationship to dominance

Examples

Cournot Competition

Cartels

	a	b
A	3, 4	4, 3
B	5, 3	3, 5
C	5, 3	4, 3

There is no strictly dominated strategy

	a	b
A	3, 4	4, 3
B	5, 3	3, 5
C	5, 3	4, 3

	a	b
A	3, 4	4, 3
B	5, 3	3, 5
C	5, 3	4, 3

- There is no strictly dominated strategy
- However, C always gives at least the same utility to player 1 as B

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A	3, 4	4, 3
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- There is no strictly dominated strategy
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- It's tempting to think player 1 would never play C

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A	3, 4	4, 3
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- There is no strictly dominated strategy
- However, C always gives at least the same utility to player 1 as B
- It's tempting to think player 1 would never play C
- However, if player 1 is sure that player two is going to play a he would be completely indifferent between playing B or C

Definition

s_i weakly dominates s'_i if for all opponent pure strategy profiles $s_{-i} \in S_{-i}$,

$$u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$$

and there is at least one opponent strategy profile $s_{-i} \in S_{-i}$ for which

$$u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i}).$$

- Given the assumptions we have, we can not eliminate a weakly dominated strategy

- Given the assumptions we have, we can not eliminate a weakly dominated strategy
- Rationality is not enough

$$\frac{dy}{ds_{100}} = -z \left(s_1 - \frac{3}{2}a - 1 \right) \left(-\frac{3}{2} \frac{1}{s_1} \right) + \lambda 100 \left(-z \left(100 - \frac{z}{2} \right) \right) - z$$

$$s_1 = s_2 = \dots = s_{100} = s$$

$$S = (z_0, z_0, \dots, z_0)$$

$$u_i(z_0, \dots, z_0) \geq u_i(60, 60, \dots, 60)$$

$$\lambda z = \frac{z \cdot 99}{-3}$$

- ▶ Given the assumptions we have, we can not eliminate a weakly dominated strategy
- ▶ Rationality is not enough
- ▶ Even so, it sounds "logical" to do so and has the potential to greatly simplify a game

- ▶ Given the assumptions we have, we can not eliminate a weakly dominated strategy
 - ▶ Rationality is not enough
 - ▶ Even so, it sounds "logical" to do so and has the potential to greatly simplify a game
- There is a problem, and that is that the order in which we eliminate the strategies matters

	a	b
A	3, 4	4, 3
C	5, 3	4, 3

$C \succ B \Rightarrow a \succ b \Rightarrow C \succ A \rightarrow \text{Sol} = (C, a)$
 $C \succ A \Rightarrow b \succ a \Rightarrow C \succ B \rightarrow \text{Sol} = (C, b)$
 $C \succ A \wedge C \succ B \Rightarrow a \wedge b \rightarrow \text{no solvable}$

- ▶ If we eliminate B (C dominates weakly), then a weakly dominates b and we can eliminate b and therefore player 1 would never play A . This leads to the result (C, a) .

	a	b
A	3, 4	4, 3
B	5, 3	3, 5
C	5, 3	4, 3

- ▶ If we eliminate B (C dominates weakly), then a weakly dominates b and we can eliminate b and therefore player 1 would never play A . This leads to the result (C, a) .
- ▶ If on the other hand, we notice that A is also weakly dominated by C then we can eliminate it in the first round, and this would eliminate a in the second round and therefore B would be eliminated. This would result in (C, b) .

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Remember the definition of competitive equilibrium in a market economy.
 Definition
 A competitive equilibrium in a market economy is a vector of prices and baskets x_i such that: 1) x_i maximizes the utility of each individual given the price vector i.e.

$$x_i = \arg \max_{x_i \in X_i, p \cdot x_i \leq p \cdot w_i} u(x_i)$$

2) the markets empty.

$$\sum_i x_i = \sum_i w_i$$

► 1) means that given the prices, individuals have no incentive to demand a different amount

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► The idea is to extend this concept to strategic situations

Best response

We denote $BR_i(s_{-i})$ (best response) as the set of strategies of individual i that maximize her utility given that other individuals follow the strategy profile s_{-i} .
 Formally,

Handwritten: \rightarrow MZ: (S_{-i})

Best response

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 Definition
 Given a strategy profile of opponents s_{-i} , we can define the best response of player i :

$$BR_i(s_{-i}) = \arg \max_{s_i \in S_i} u(s_i, s_{-i})$$

► $s_i \in BR_i(s_{-i})$ if and only if $u(s_i, s_{-i}) \geq u(s'_i, s_{-i})$ for all $s'_i \in S_i$.

Handwritten:

$$\max_x f(x) = f(x^*)$$

$$\text{ARG MAX}_x f(x) = x^*$$

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► $s_i \in BR_i(s_{-i})$ if and only if $u(s_i, s_{-i}) \geq u(s'_i, s_{-i})$ for all $s'_i \in S_i$.

► There could be multiple strategies in $BR_i(s_{-i})$ but all such strategies give the same utility to player i if the opponents are indeed playing according to s_{-i} .

Nash equilibrium

Normal

Definition
 Suppose that we have a game $(I = \{1, 2, \dots, n\}, S_1, \dots, S_n, u_1, \dots, u_n)$. Then a strategy profile $s^* = (s_1^*, \dots, s_n^*)$ is a **pure strategy Nash equilibrium** if for every i and for every $s_i \in S_i$,

$$u_i(s_i^*, s_{-i}^*) \geq u_i(s_i, s_{-i}^*)$$

$s_i^* \in MR_i(s_{-i}^*)$

◀ ▶ ↻ 🔍

Nash equilibrium

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- ▶ Analogous to that of a competitive equilibrium in the sense that nobody has unilateral incentives to deviate

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- ▶ once this equilibrium is reached, nobody has incentives to move from there

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- ▶ Analogous to that of a competitive equilibrium in the sense that nobody has unilateral incentives to deviate
- ▶ once this equilibrium is reached, nobody has incentives to move from there
- ▶ This is a concept of stability, but there is no way to ensure, or predict, that the game will reach this equilibrium

◀ ▶ ↻ 🔍

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Beauty contest

► Consider the following game among 2 people. Each individual selects a number, s_i , between 20 and 60.

Beauty contest

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► Let s_{-i} be the number selected by the other individual.

Beauty contest

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► Let s_{-i} be the number selected by the other individual.

► The utility function of the individual i is $U_i(s_i, s_{-i}) = 100 - (s_i - \frac{2}{3}s_{-i})^2$

$U_1(s_1, s_2) = 100 - (s_1 - \frac{2}{3}s_2)^2$
 $U_2(s_1, s_2) = 100 - (s_2 - \frac{2}{3}s_1)^2$

Beauty contest

The best response of an individual is given by

$$s_i(s_{-i}) = \begin{cases} \frac{3}{2}s_{-i} & \text{if } s_{-i} \leq 40 \\ 60 & \text{if } s_{-i} > 40 \end{cases}$$

The Nash equilibrium is where both BR functions intersect (i.e., when both play 60)

Prisoner's dilemma

	s_2	C	NC
s_1	C	5,5	0,10
	NC	10,0	2,2

$MR_1(s_2) = NC$
 $MR_2(s_1) = NC$

$EN = (NC, NC)$

Prisoner's dilemma

	C	NC
C	5,5	0,10
NC	10,0	2,2

The best response functions are:

$$BR_i(s_{-i}) = \begin{cases} NC & \text{if } s_{-i} = C \\ C & \text{if } s_{-i} = NC \end{cases}$$

The Nash equilibrium is where both BR functions intersect (i.e., when both play NC, i.e., (NC, NC))

$MR_i(s_i) \Rightarrow \text{OPTIMA}$

$$\frac{\partial U(s_i, s_{-i})}{\partial s_i} = -2(s_i - \frac{2}{3}s_{-i}) = 0$$

$$s_i = \frac{2}{3}s_{-i} \Rightarrow s_i = \begin{cases} \frac{3}{2}s_{-i} & \text{if } s_{-i} \in [20, 40] \\ 60 & \text{if } s_{-i} > 40 \end{cases}$$

$U_i(s_i, s_{-i}) = 100 - (s_i - \frac{2}{3}s_{-i})^2$

$s^* = (60, 60)$

$s_i^* = \frac{3}{2}s_{-i}$
 $s_{-i}^* = \frac{3}{2}s_i$

$U_i(20, 20) > U(60, 60)$

$U_1(30, 20) > U_1(20, 20) \rightarrow$ SI INCENTIVOS UNICATERALES A FORTALESE

$U_i(s_i, 60) \leq U_i(60, 60)$

Prisoner's dilemma – A trick

Best response of 1 to 2 playing C

	C	NC
C	5,5	0,10
NC	10,0	2,2

Handwritten notes: $S = \{(NC, NC)\}$, $E = \{C\}$

Prisoner's dilemma – A trick

Best response of 1 to 2 playing NC

	C	NC
C	5,5	0,10
NC	10,0	2,2

Prisoner's dilemma – A trick

Best response of 2 to 1 playing C

	C	NC
C	5,5	0,10
NC	10,0	2,2

Prisoner's dilemma – A trick

Best response of 2 to 1 playing NC

	C	NC
C	5,5	0,10
NC	10,0	2,2

When underlined for both players, it is a Nash equilibrium (both are doing their BR)

Battle of the sexes

	G	P
G	2,1	0,0
P	0,0	1,2

Handwritten notes: $S = \{(G,G), (P,P)\}$

Battle of the sexes

	G	P
G	2,1	0,0
P	0,0	1,2

$$BR(s_{-i}) = \begin{cases} G & \text{if } s_{-i} = G \\ P & \text{if } s_{-i} = P \end{cases}$$

Battle of the sexes

	G	P
G	2,1	0,0
P	0,0	1,2

$$BR_i(s_{-i}) = \begin{cases} G & \text{if } s_{-i} = G \\ P & \text{if } s_{-i} = P \end{cases}$$

Thus, (G, G) y (P, P) are both Nash equilibrium

Matching pennies (Pares o Nones) – Simultaneous

S_2

S_1

	1	2
1	(1000,-1000)	(-1000,1000)
2	(-1000,1000)	(1000,-1000)

Matching pennies (Pares o Nones) – Simultaneous

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Matching pennies (Pares o Nones) – Simultaneous

	1	2
1	(1000,-1000)	(-1000,1000)
2	(-1000,1000)	(1000,-1000)

$$BR_1(s_2) = \begin{cases} 1 & \text{if } s_2 = 1 \\ 2 & \text{if } s_2 = 2 \end{cases}$$

$$BR_2(s_1) = \begin{cases} 2 & \text{if } s_1 = 1 \\ 1 & \text{if } s_1 = 2 \end{cases}$$

There is no Nash equilibrium in pure strategies

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Dominance

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Nash equilibrium survive IDSDS

Theorem:
Every Nash equilibrium survives the iterative elimination of strictly dominated strategies

↳ No es eliminado

Proof:
By contradiction:
▶ Suppose it is not true

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▶ Without loss of generality say we eliminated the strategy s_i^* of individual i

Proof:
By contradiction:
▶ Suppose it is not true
▶ Then we must have eliminated some strategy in the Nash equilibrium s^*
▶ Lets zoom in in the round where we first eliminate a strategy that is part of s^*
▶ Without loss of generality say we eliminated the strategy s_i^* of individual i
▶ It must have been that
$$u_i(s_i^*, s_{-i}) < u_i(s_i, s_{-i}) \quad \forall s_{-i} \in S_{-i}$$

Proof

By contradiction:

- Suppose it is not true
- Then we must have eliminated some strategy in the Nash equilibrium s^*
- Lets zoom in in the round where we first eliminate a strategy that is part of s^*
- Without loss of generality say we eliminated the strategy s_i^* of individual i
- It must have been that

$$u(s_i^*, s_{-i}) < u(s_i, s_{-i}) \forall s_{-i} \in S_{-i}$$

*Si Damsam A Si**

...

- In particular

$$u(s_i^*, s_{-i}^*) < u(s_i, s_{-i}^*)$$

$\Rightarrow S_i^* \& \text{MIZI } (S_i^*)$

Proof

By contradiction:

- Suppose it is not true
- Then we must have eliminated some strategy in the Nash equilibrium s^*
- Lets zoom in in the round where we first eliminate a strategy that is part of s^*
- Without loss of generality say we eliminated the strategy s_i^* of individual i
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- In particular

$$u(s_i^*, s_{-i}^*) < u(s_i, s_{-i}^*)$$

- But this means s_i^* is not the best response of individual i to s_{-i}^*

Proof

By contradiction:

- Suppose it is not true
- Then we must have eliminated some strategy in the Nash equilibrium s^*
- Lets zoom in in the round where we first eliminate a strategy that is part of s^*
- Without loss of generality say we eliminated the strategy s_i^* of individual i
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...

- In particular

$$u(s_i^*, s_{-i}^*) < u(s_i, s_{-i}^*)$$

- But this means s_i^* is not the best response of individual i to s_{-i}^*
- And this is a contradiction!

Nash equilibrium survive IDSDS

Theorem

If the process of IDSDS comes to a single solution, that solution is a Nash Equilibrium and is unique.

↳ Successes Soluble Part I IDSDS

Proof

First let's proof its a Nash Equilibrium. The fact that is unique is trivial by the previous theorem.

Proof:

By contradiction:

- Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium

□

Proof

First let's proof its a Nash Equilibrium. The fact that is unique is trivial by the previous theorem.

Proof:

By contradiction:

- Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium
- For some individual i there exists s_i such that

$$u(s_i, s_{-i}^*) > u(s_i^*, s_{-i}^*)$$

□

Proof

First let's prove it's a Nash Equilibrium. The fact that it is unique is trivial by the previous theorem.

Proof:

By contradiction:

- Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium
- For some individual i there exists s_i such that

$$u_i(s_i, s_{-i}^*) > u_i(s_i^*, s_{-i}^*)$$
- But then s_i could not have been eliminated

□

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50

Proof

First let's prove it's a Nash Equilibrium. The fact that it is unique is trivial by the previous theorem.

Proof:

By contradiction:

- Suppose that the results from IDSDS (s^*) is not a Nash Equilibrium
- For some individual i there exists s_i such that

$$u_i(s_i, s_{-i}^*) > u_i(s_i^*, s_{-i}^*)$$
- But then s_i could not have been eliminated
- And this is a contradiction!

□

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 - Cournot Competition

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50

Cournot Competition

- We will apply the concept of pure Nash equilibrium to analyze oligopoly markets

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Cournot Competition

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Cournot Competition → **COMPETEN EN CANTIDADES**

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$$P(Q) = 120 - Q, Q = q_1 + q_2.$$
- ▶ Strategy space is $S_i = [0, +\infty)$ → **CANTIDADES.**
- ▶ The utility function of player i is given by:

$$\pi_1(q_1, q_2) = (120 - (q_1 + q_2))q_1$$

$$\pi_2(q_1, q_2) = (120 - (q_1 + q_2))q_2$$

Handwritten notes:

$\frac{\partial \pi_1}{\partial q_1} = 120 - 2q_1 - q_2 \Rightarrow \frac{120 - q_2}{2} = q_1$

$\pi(q_1, q_2) = P(Q) \cdot q_1 \Rightarrow 0$

Cournot Competition

- ▶ Are there any strictly dominant strategies?

Cournot Competition

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Cournot Competition

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- ▶ Are there any others? given q_{-i} ,

$$M_i \Rightarrow \frac{d\pi_i}{dq_i}(120 - q_i - q_{-i})q_i = 120 - 2q_i - q_{-i} \rightarrow \frac{120 - q_{-i} - q_i}{2}$$

Cournot Competition

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$$\frac{d\pi_i}{dq_i}(120 - q_i - q_{-i})q_i = 120 - 2q_i - q_{-i}$$

- ▶ Therefore 60 strictly dominates any $q_i \in (60, 120] \rightarrow q_i, q_{-i} \in [0, 60]$

Cournot Competition

- ▶
$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$$

Cournot Competition

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$$BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$$
- ▶ for any $q_i \in [0, 60]$, there exists some $q_{-i} \in [0, +\infty)$ such that $BR_i(q_{-i}) = q_i$

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ for any $q_i \in [0, 60]$, there exists some $q_{-i} \in [0, +\infty)$ such that $BR_i(q_{-i}) = q_i$
- ▶ Such a q_i can never be strictly dominated

Navigation icons

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
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- ▶ After one round of deletion of strictly dominated strategies, we are left with: $S_i = [0, 60]$

Navigation icons

Cournot Competition

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Navigation icons

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ $q_{-i} = [0, 60]$

Navigation icons

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ $q_{-i} = [0, 60]$
- ▶ Therefore $q_i \in [0, 30]$ are strictly dominated by $q_i = 30$

Navigation icons

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ $q_{-i} = [0, 60]$
- ▶ Therefore $q_i \in [0, 30]$ are strictly dominated by $q_i = 30$
- ▶ After two rounds of deletion of strictly dominated strategies, we are left with: $S_i = [30, 60]$

Navigation icons

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ $q_{-i} = [30, 60]$
- ▶ 45 strictly dominates all strategies $q_i \in (45, 60]$
- ▶ After three rounds of deletion of strictly dominated strategies, we are left with: $S_i = [30, 45]$

Cournot Competition

- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ $q_{-i} = [30, 45]$
- ▶ 37.5 strictly dominates all strategies $q_i \in [30, 37.5]$
- ▶ After four rounds of deletion of strictly dominated strategies, we are left with: $S_i = [37.5, 45]$

Cournot Competition

- ▶ After (infinitely) many iterations, the only remaining strategies are $S_i = 40$
- ▶ The unique solution by IDSDS is $q_1^* = q_2^* = 40$. $\rightarrow \text{EN}$

Cournot Competition

- ▶ There will also be a unique Nash equilibrium

Cournot Competition

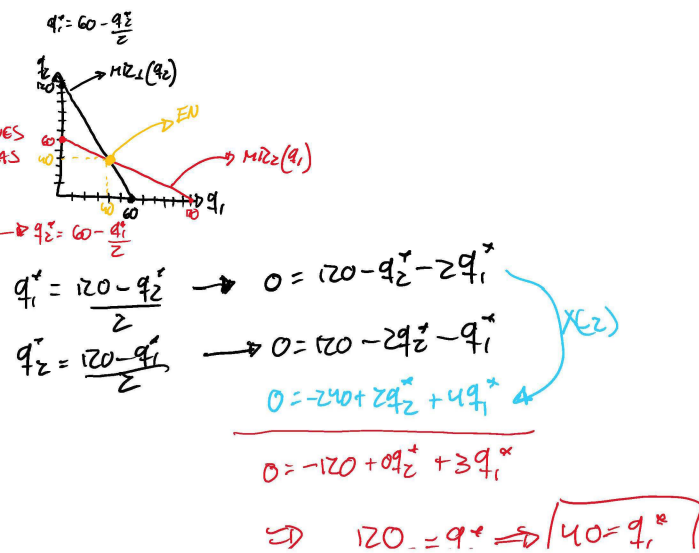
- ▶ There will also be a unique Nash equilibrium
- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2} = MR_i(q_{-i})$
- ▶ $q_1^* = \frac{120 - q_2^*}{2} \Rightarrow 2 \text{ ECUACIONES}$
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Cournot Competition

- ▶ There will also be a unique Nash equilibrium
- ▶ $BR_i(q_{-i}) = \frac{120 - q_{-i}}{2}$
- ▶ At any Nash equilibrium, we must have: $q_1^* \in BR_1(q_2^*)$ and $q_2^* \in BR_2(q_1^*)$.

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$$q_1^* = \frac{120 - q_2^*}{2}, q_2^* = \frac{120 - q_1^*}{2}$$

Cournot Competition

- There will also be a unique Nash equilibrium
- At any Nash equilibrium, we must have: $q_1^* \in BR_1(q_2^*)$ and $q_2^* \in BR_2(q_1^*)$.
- We can solve for q_1^* and q_2^* to obtain:

$$q_1^* = 40, q_2^* = 40, Q^* = 80, \Pi_1^* = \Pi_2^* = 1600$$

Cournot Competition vs Monopoly (cartel)

- In a perfectly competitive market, price equals marginal cost and the total quantity produced will be $Q = 120$.

Cournot Competition vs Monopoly (cartel)

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- A monopolist would solve the following maximization problem:

$$\max_Q (120 - Q)Q \rightarrow Q^* = 60, P^* = 60, \Pi^* = 3600$$

$$\pi = (120 - Q)Q$$

$$\frac{\partial \pi}{\partial Q} = 120 - 2Q = 0 \rightarrow Q = 60$$

$$P = 60$$

Cournot Competition vs Monopoly (cartel)

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- The profits to each firm in the Cournot Competition is less than half of the monopoly profits

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- The profits to each firm in the Cournot Competition is less than half of the monopoly profits
- In a duopoly, externalities are imposed on the other firm

Cournot Competition - General case

- n firms are competing a la Cournot

$$0 = -120 + 0q_2 + 3q_1$$

$$\Rightarrow \frac{120}{3} = q_1^* \Rightarrow 40 = q_1^*$$

$$q_2^* = \frac{120 - q_1^*}{2} = \frac{120 - 40}{2} = \frac{80}{2} = 40$$

$$Q_1 = Q_2 = 30 \text{ i. } Q = 60 = Q^M$$

$$\Pi_1(q_1, q_2 = 30) = (120 - q_1 - 30)q_1$$

$$\frac{\partial \Pi}{\partial q_1} = 90 - 2q_1 = 0$$

$$q_1^* = 45$$

$$\sum q_i = 25$$

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- To simplify notation, let $Q_{-i} = \sum_{j \neq i} q_j$
- $$\max_{q_i} p(q_i + Q_{-i})q_i - c_i(q_i)$$

Cournot Competition - General case

$$\max_{q_i} p(q_i + Q_{-i})q_i - c_i(q_i)$$

- First order condition implies:

$$q_i \frac{dp}{dq}(q_i + Q_{-i}) + P(q_i + Q_{-i}) = \frac{dc_i}{dq}(q_i)$$

$$q_i \frac{dp}{dq}(Q) + P(Q) = \frac{dc_i}{dq}(q_i)$$

$$P(Q) - \frac{dc_i}{dq}(q_i) = -q_i \frac{dp}{dq}(Q)$$

$$\frac{P(Q) - \frac{dc_i}{dq}(q_i)}{P(Q)} = -\frac{q_i}{Q} \frac{dp}{dq}(Q)$$

$$\frac{P(Q) - \frac{dc_i}{dq}(q_i)}{P(Q)} = -\frac{q_i}{Q} \epsilon_{Q,p}(Q)$$

Cournot Competition - General case

$$\frac{P(Q) - \frac{\partial c_i}{\partial q_i}(q_i)}{P(Q)} = \frac{q_i}{Q} \frac{1}{\varepsilon_{Q,P}(Q)}$$

► Therefore in a pure strategy Nash equilibrium $(q_1^*, q_2^*, \dots, q_n^*)$ with $Q^* = q_1^* + q_2^* + \dots + q_n^*$, we must have:

$$\frac{P(Q^*) - \frac{\partial c_1}{\partial q_1}(q_1^*)}{P(Q^*)} = \frac{q_1^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}$$

$$\frac{P(Q^*) - \frac{\partial c_2}{\partial q_2}(q_2^*)}{P(Q^*)} = \frac{q_2^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}$$

$$\vdots$$

$$\frac{P(Q^*) - \frac{\partial c_n}{\partial q_n}(q_n^*)}{P(Q^*)} = \frac{q_n^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}$$

Cournot Competition - General case

► Suppose that all firms have exactly the same cost function c

$$\frac{P(Q^*) - \frac{\partial c}{\partial q_i}(q_i^*)}{P(Q^*)} = \frac{q_i^*}{Q^*} \frac{1}{\varepsilon_{Q,P}(Q^*)}$$

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► Let us conjecture that there exists a pure strategy Nash equilibrium that is **symmetric**, in which $q_1^* = q_2^* = \dots = q_n^* = q^*$

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► In this case $Q^* = nq^*$

$$\frac{P(nq^*) - \frac{\partial c}{\partial q_i}(q^*)}{P(nq^*)} = \frac{1}{n} \frac{1}{\varepsilon_{Q,P}(nq^*)}$$

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► Rewriting

$$P(Q^*) = \frac{1}{1 + \frac{1}{n} \frac{1}{\varepsilon_{Q,P}(Q^*)}} \frac{\partial c}{\partial q} \left(\frac{Q^*}{n} \right)$$

Lecture 12: Game Theory // Nash equilibrium

- Dominance
 - Weakly dominated strategies
- Nash equilibrium
- Some examples
- Relationship to dominance
- Examples
 - Cournot Competition
 - Cartels

Cartels

- Suppose there are three firms who face zero marginal cost
- The inverse demand function is given by:

$$p(q_1 + q_2 + q_3) = 1 - q_1 - q_2 - q_3 = 1 - Q$$

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$$1 - 2q_j - Q_{-j} = 0 \implies q_j = \frac{1 - Q_{-j}}{2} \implies BR_j(Q_{-j}) = \frac{1 - Q_{-j}}{2}$$

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- In a Nash equilibrium we must have:

$$\begin{aligned} q_1^* &= \frac{1 - q_2^* - q_3^*}{2} \\ q_2^* &= \frac{1 - q_1^* - q_3^*}{2} \\ q_3^* &= \frac{1 - q_1^* - q_2^*}{2} \end{aligned}$$

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Cartels

- The easiest way to solve this first, let us add the three equations to get:

$$Q^* = \frac{3}{2} - Q^* \implies Q^* = \frac{3}{4}$$

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- $q_1^* = q_2^* = q_3^* = \frac{1}{4}$

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- ▶ $q_1^* = q_2^* = q_3^* = \frac{1}{4}$
- ▶ Price is $p^* = 1/4$ and all firms get the same profits of $1/16$

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Cartels

- ▶ Two of the firms merge into firm A, while one of the firms remains single, call that firm B

◀ ▶ ↻ 🔍

Cartels

- ▶ Two of the firms merge into firm A, while one of the firms remains single, call that firm B
- ▶ Each firm then again faces the profit maximization problem:

$$\max_q (1 - q - q_{-i})q \implies BR_i(q_{-i}) = \frac{1 - q_{-i}}{2}$$

◀ ▶ ↻ 🔍

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- ▶ Therefore

$$q_A^* = \frac{1 - q_B^*}{2}$$

$$q_B^* = \frac{1 - q_A^*}{2}$$

◀ ▶ ↻ 🔍

Cartels

- ▶ Solving this:

$$q_A^* = q_B^* = \frac{1}{3}$$

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Cartels

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- ▶ The price is then $p^* = 1/3$

◀ ▶ ↻ 🔍

Cartels

- ▶ Solving this:

$$q_A^c = q_B^c = \frac{1}{3}$$

- ▶ The price is then $p^* = 1/3$
- ▶ If the profits are shared equally among firms 1 and 2 who have merged, then profits of firms 1 and 2 are $1/18$ whereas firm 3 obtains a profit of $1/9$

Cartels

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- ▶ If the profits are shared equally among firms 1 and 2 who have merged, then profits of firms 1 and 2 are $1/18$ whereas firm 3 obtains a profit of $1/9$
- ▶ Firms 1 and 2 suffered, while firm 3 is better off!
- ▶ Firm 3 is obtaining a disproportionate share of the joint profits (more than $1/3$)

Cartels

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- ▶ Total profits then are given by $\frac{1}{4}$ which means that each firm obtains a profit of $\frac{1}{12} < \frac{1}{9}$

Cartels

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► In the monopolist problem, we solve:

$$\max_Q (1 - Q)Q \Rightarrow Q^* = \frac{1}{2}$$

► Total profits then are given by $\frac{1}{4}$ which means that each firm obtains a profit of $\frac{1}{8} < \frac{1}{4}$

► Firm 3 clearly wants to stay out

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Cartels

There are many difficulties associated with sustaining collusive agreements (e.g., the OPEC cartel)

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