



Lecture 17: Applications of Subgame Perfect Nash Equilibrium

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- We will use (G, T) to denote that game G is repeated T times

LOS ANGELES BASE

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 2. Players observe the actions chosen by the players in period 1. Then in period 2, players simultaneously play the game G .
 3. This game proceeds until time T .
 4. After time T , if the action profiles chosen in times $1, 2, \dots, T$ are given by $((a_1^1, a_2^1), \dots, (a_1^T, a_2^T))$,

$$\sum_{t=1}^T \beta^{t-1} v(a_1^t, a_2^t) = \text{Valor Presente} = \sum \text{Pagos}$$

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- Working incurs a cost of 1 however increases the utility of the other player $-i$ by 2
- Thus, $u_i(e_1, e_2) = 2e_{-i} - e_i$.

Prisoner's Dilemma (Game G)

| | | |
|-----------|-----------|-----------|
| | $e_2 = 1$ | $e_2 = 0$ |
| $e_1 = 1$ | 1, 1 | 1, 0 |
| $e_1 = 0$ | 0, 1 | 0, 0 |

$EU = (e_1=0, e_2=0)$

- What happens when $T = 1$

► What happens when $T = 1$

► NE: Players 1 and 2 will both choose $(e_1 = 0, e_2 = 0)$

Imagine players are engaged in a long run relationship that lasts more than just playing the game once: (G, Z)

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- Both players play the simultaneous move game G .
- Both players observe the actions chosen by the two players. Then they play G again.
- Then payoffs are realized as the discounted sum of the utilities of the actions in each period with discount factor $\delta \in (0, 1)$.

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$u_1 = 1 + \delta \cdot 2$
 $u_2 = 1 + \delta \cdot (-1)$

► We will solve for the set of pure SPNE of this game.

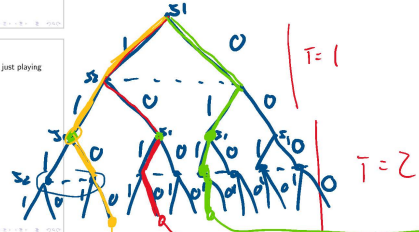
► Player 1 has 5 information sets in total

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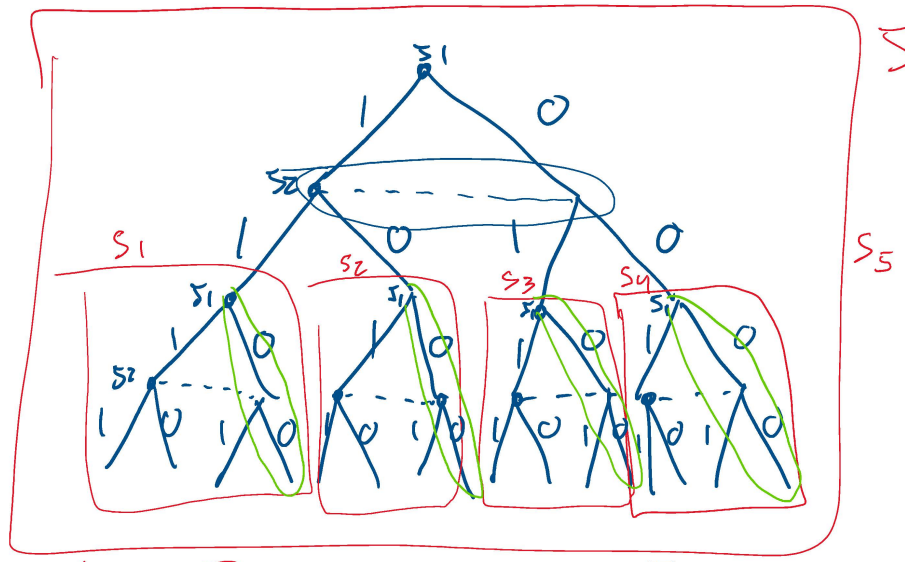
► Player 1 has 5 information sets in total

► A pure strategy for player 1 must specify what he does in each of these information sets

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$u_1 = 1 + \delta$
 $u_2 = 1 + \delta$
 $u_1 = -1 + (-1)\delta$
 $u_2 = 2 + (2)\delta$
 $u_1 = 2 + 1\delta$
 $u_2 = -1 + 1\delta$



S SUBSUBGOS

$S_1 - S_4 \rightarrow$

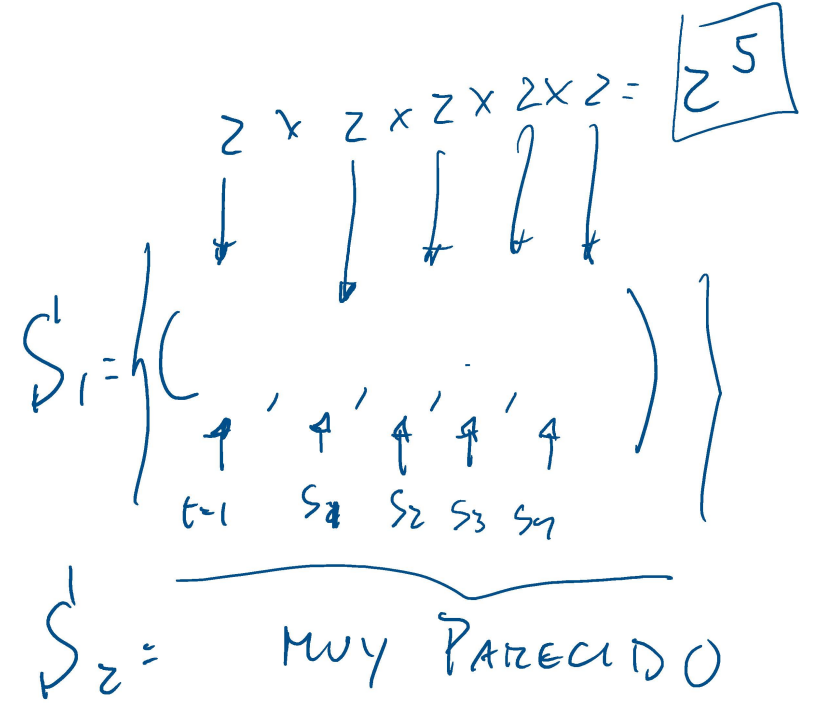
| | | | |
|-------|---|--|--|
| | | 1 | 0 |
| S_1 | 1 | $U_{t=1}^1 + 1\delta, U_{t=1}^2 + 1\delta$ | $U_{t=1}^1 - 1\delta, U_{t=1}^2 + 2\delta$ |
| | 0 | $U_{t=1}^1 + 2\delta, U_{t=1}^2 - 1\delta$ | $U_{t=1}^1 + 0\delta, U_{t=1}^2 + 0\delta$ |

$(U_{t=1}^1, U_{t=1}^2) +$

| | | | |
|---|---|---------------------|---------------------|
| | | 1 | 0 |
| 1 | 1 | $1\delta, 1\delta$ | $-1\delta, 2\delta$ |
| | 0 | $2\delta, -1\delta$ | $0\delta, 0\delta$ |

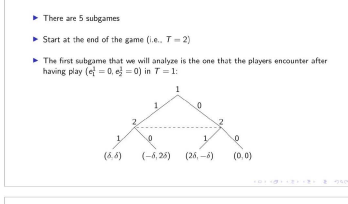
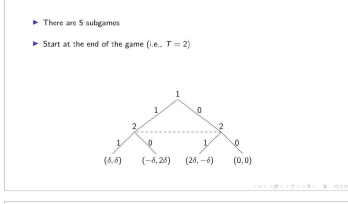
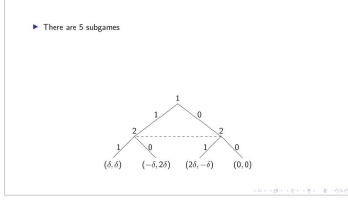
$(U_{t=1}^1, U_{t=1}^2) + \delta$

| | | | |
|---|---|---------|---------|
| | | 1 | 0 |
| 1 | 1 | $1, 1$ | $-1, 2$ |
| | 0 | $2, -1$ | $0, 0$ |



► We will solve for the set of pure SPNE of this game.
 ► Player 1 has 5 information sets in total.
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 ► Player 1 has a total of $32 (2^5)$ pure strategies.

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 ► Player 1 has 5 information sets in total.
 ► A pure strategy for player 1 must specify what he does in each of these information sets.
 ► Player 1 has a total of $32 (2^5)$ pure strategies.
 ► Similarly, player 2 has a total of 32 pure strategies.



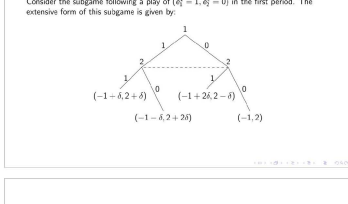
The Nash equilibria can be seen by writing out the normal form of the game.

Normal Form of Extensive Form

| | | |
|-----------|---------------|---------------|
| | $e_2 = 1$ | $e_2 = 0$ |
| $e_1 = 1$ | $6, 6$ | $-6, 2\delta$ |
| $e_1 = 0$ | $2\delta, -6$ | $0, 0$ |

► This game has a unique Nash equilibrium in which the players play $(e_1^1 = 0, e_2^1 = 0)$

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 ► Therefore after having observed $(e_1^1 = 0, e_2^1 = 0)$ in the first period, both players will play $(e_1^2 = 0, e_2^2 = 0)$ in period 2



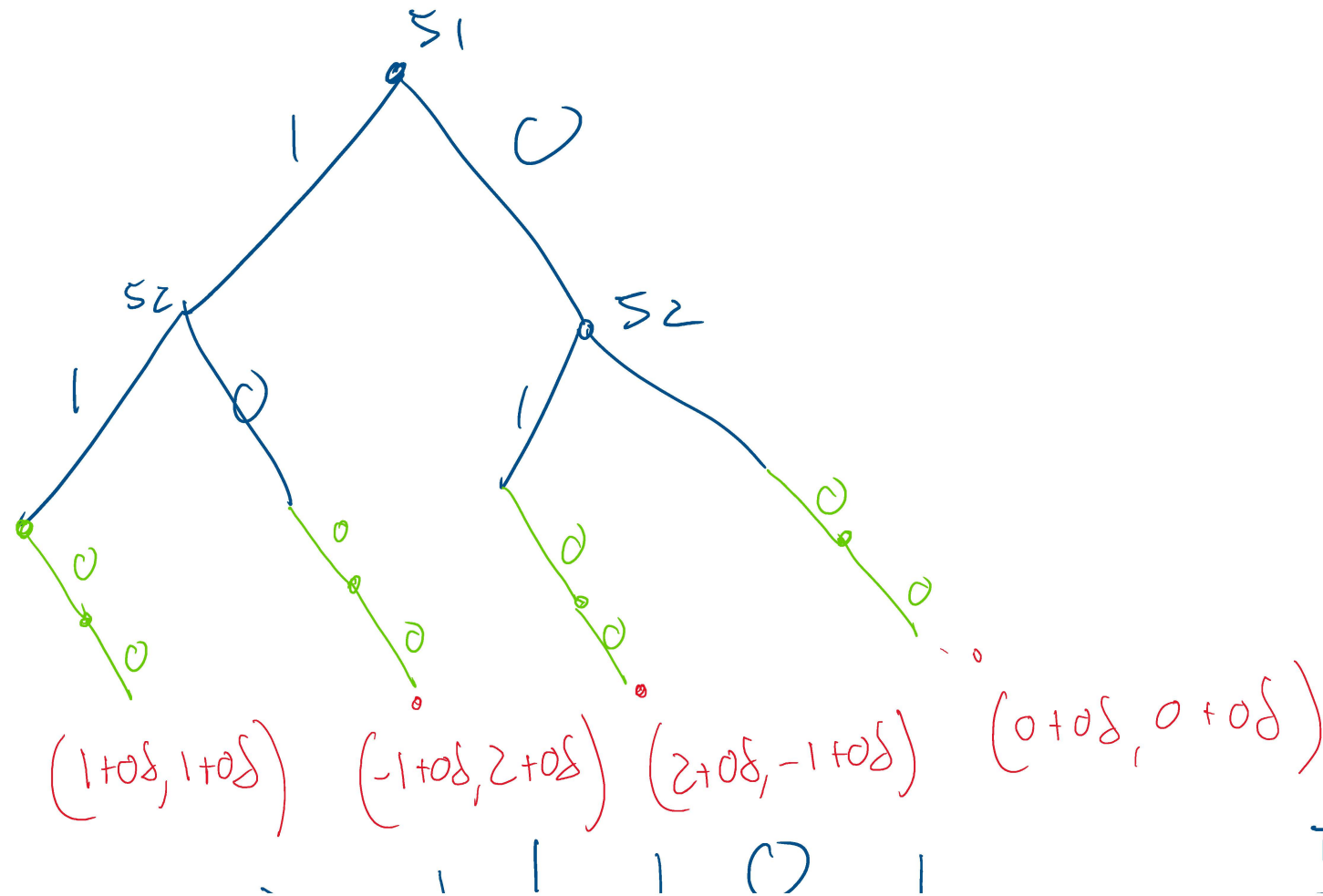
$$(V_{e_1=1}, V_{e_2=1}) + \delta \begin{array}{c|cc} & 1 & 0 \\ \hline 1 & 1, 1 & -1, 2 \\ \hline 0 & 2, -1 & 0, 0 \end{array}$$

$$(V_{e_1=1}, V_{e_2=1}) + \delta \begin{array}{c|cc} & 1 & 0 \\ \hline 1 & 1, 1 & -1, 2 \\ \hline 0 & 2, -1 & 0, 0 \end{array}$$

MISSTOS
 INCENTIVOS
 DE G

$$\exists N_{\text{SUBJUGO}} = (e_1 = 0, e_2 = 0)$$

$T=1$ \rightarrow S_5 subgame completo



PAGO EN $T=2$

$T=1$

$(-1, 2)$ $(-1, 2)$

The normal form of this subgame can be seen in the Table

| | $e_2 = 1$ | $e_2 = 0$ |
|-----------|-------------------------|--------------------------|
| $e_1 = 1$ | $-1 + 2\delta + \delta$ | $-1 - 2\delta + 2\delta$ |
| $e_1 = 0$ | $-1 + 2\delta - \delta$ | $-1, 2$ |

► $(e_1 = 0, e_2 = 0)$ is the unique Nash equilibrium

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► In any SPNE, $(e_1^1 = 0, e_2^1 = 0)$ must be played after observing $(e_1^1 = 1, e_2^1 = 0)$

► We can go through the remaining smaller subgames after the observation of $(e_1^1 = 1, e_2^1 = 0)$ and after the observation of $(e_1^1 = 1, e_2^1 = 1)$

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► We will reach the same conclusion in each of these scenarios: that $(e_1^1 = 0, e_2^1 = 0)$ must be played in each of these subgames

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► Regardless of the observed action, $(0, 0)$ is played in period 2

► Why is this the case?

► The idea is that payoffs that have accrued in period 1 are essentially sunk, and have no influence on incentives in period 2

To see this consider the normal form representation in the subgame after the observation of $(e_1^1 = 1, e_2^1 = 0)$

| | $e_2 = 1$ | $e_2 = 0$ |
|-----------|-------------------------|--------------------------|
| $e_1 = 1$ | $-1 + 2\delta + \delta$ | $-1 - 2\delta + 2\delta$ |
| $e_1 = 0$ | $-1 + 2\delta - \delta$ | $-1, 2$ |

S_1

| | | |
|---|-----------------------------|-----------------------------|
| | 1 | 0 |
| 1 | $1 + 0\delta, 1 + 0\delta$ | $-1 + 0\delta, 2 + 0\delta$ |
| 0 | $2 + 0\delta, -1 + 0\delta$ | $0 + 0\delta, 0 + 0\delta$ |

$T=1$

| | |
|-------|-------|
| 1, 1 | -1, 2 |
| 2, -1 | 0, 0 |

+ $(0, 0)\delta$

$EN_{SS} (e_1 = 0, e_2 = 0)$

MOMOS
INCENTIVOS
JUEGO
BASE

EPS = $(S_1 = (0, 0, 0, 0, 0), S_2 = (0, 0, 0, 0, 0))$

- We can subtract off the payoff that player 1 received in period 1 and divide through player 1's payoffs by 2 to obtain the following payoff matrix

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| | $a_1=1$ | $a_1=0$ |
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- We can subtract off the payoff that player 1 received in period 1 and divide through player 1's payoffs by 2 to obtain the following payoff matrix

- We can do the same thing for player 2's payoffs and get the payoff matrix

Normal Form of Extensive Form

| | $a_2=1$ | $a_2=0$ |
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| $a_1=1$ | 1, 1 | -1, 2 |
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- This will be true no matter the action profile played in period 1

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► So what have we learned?

- Basically after any history, the strategic normal form is essentially the same as the original prisoner's dilemma
- Both players play ($e_1^t = 0, e_2^t = 0$) after any information set in the last period

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► Now let us see what must be played in the first period by the two players

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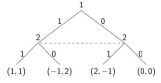
- Both players anticipate that ($e_1^2 = 0, e_2^2 = 0$) will be played after any chosen action profile in the first period

3 / 11

► Now let us see what must be played in the first period by the two players

- Both players anticipate that ($e_1^2 = 0, e_2^2 = 0$) will be played after any chosen action profile in the first period

► We can simplify the extensive form game to the following:



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If we draw the normal form of this game, then we get:

Normal Form of Extensive Form

| | $e_2 = 1$ | $e_2 = 0$ |
|-----------|-----------|-----------|
| $e_1 = 1$ | 1,1 | -1,2 |
| $e_1 = 0$ | 2,-1 | 0,0 |

The unique Nash equilibrium of the above normal form game is ($e_1^1 = 0, e_2^1 = 0$)

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Therefore the unique SPNE is:

$$\left(\left(\begin{matrix} e_1^1 = 0 \\ e_2^1 = 0 \\ e_1^2 = 0 \\ e_2^2 = 0 \end{matrix} \right), \left(\begin{matrix} e_1^2 = 0 \\ e_2^2 = 0 \\ e_1^1 = 0 \\ e_2^1 = 0 \end{matrix} \right) \right)$$

In other words both players always shirk

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EPS

- Here the unique **SPNE** requires all players to play $e_i = 0$ at all periods and all information sets.
- Thus, the equilibrium outcome is simply the repetition of the unique NE of the stage game.
- This holds more generally when the stage game has a **unique NE**.
- Whenever the stage game has a unique NE, then the only SPNE of a finite horizon repeated game with that stage game is the repetition of the stage game NE.

(G,T)
↳ FINITE

Theorem

Suppose that the stage game G has exactly one NE $(s_1^*, s_2^*, \dots, s_n^*)$. Then for any $\delta \in (0, 1]$ and any T , the T -stage repeated game has a unique SPNE in which all players i play s_i^* at all information sets.

en todos los periodos

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- But then we can induct.
- Knowing that the stage game Nash equilibrium is going to be played tomorrow, at any information set, we can ignore the past payoffs.
- We concentrate just on the payoffs in the future. Thus in period $T-1$, player i simply wants to maximize:

$$\max_{s_i} \delta^{T-2} u_i(s_i, s_{-i}^{T-1}) + \delta^{T-1} u_i(s_i^*)$$

G, Repetition



► What player i plays today has no consequences for what happens in period T since we saw that all players will play a^* no matter what happens in period $T - 1$

1 / 20

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► So, the maximization problem above is the same as:

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2 / 20

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► Following exactly this induction, we can conclude that every player must play a_i^* at all times and all histories

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