



Lecture17...

Lecture 17: Applications of Subgame Perfect Nash Equilibrium

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- ▶ We will use (G, T) to denote that game G is repeated T times

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3. The game proceeds until time T.
4. After time T, the utility profiles chosen in times 1, 2, ..., T are given by $(\theta_1^T, \theta_2^T, \dots, \theta_n^T)$. $\sum_{t=1}^T u_i(\theta^t, x^t)$

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- Working incurs a cost of 1, however increases the utility of the other player $-i$ by 2
- This: $u_i(a_i, a_{-i}) = 2a_{-i} - a_i$

Player's Utilities (Game G)

$a_1 \backslash a_2$	0	1
0	0	0
1	1	0

- What happens when $T = 1$

► What happens when $T = 1$

► NE: Players 1 and 2 will both choose $(e_1 = 0, e_2 = 0)$

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1. Both players play the simultaneous move game G .

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3. Then payoffs are realized as the discounted sum of the utilities of the actions in each period with discount factor $\delta \in [0, 1]$.

$U_i(e_i, e_{-i}) = z e_i - e_i$

Suppose that the two players chose $(e_1 = 1, e_2 = 1)$ in the first period
In the second period, they chose $(e_1 = 0, e_2 = 1)$

$U_1 = 1 + z\delta$
 $U_2 = 1 - \delta$

Factor de Descuento δ

Suppose that the two players chose $(e_1 = 1, e_2 = 1)$ in the first period
In the second period, they chose $(e_1 = 0, e_2 = 1)$

$u_1 = 1 + \delta - z$
 $u_2 = 1 + \delta - (-1)$

► We will solve for the set of pure SPNE of this game.

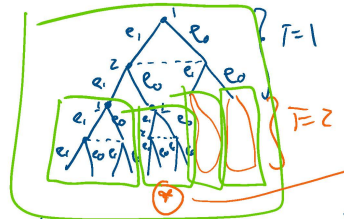
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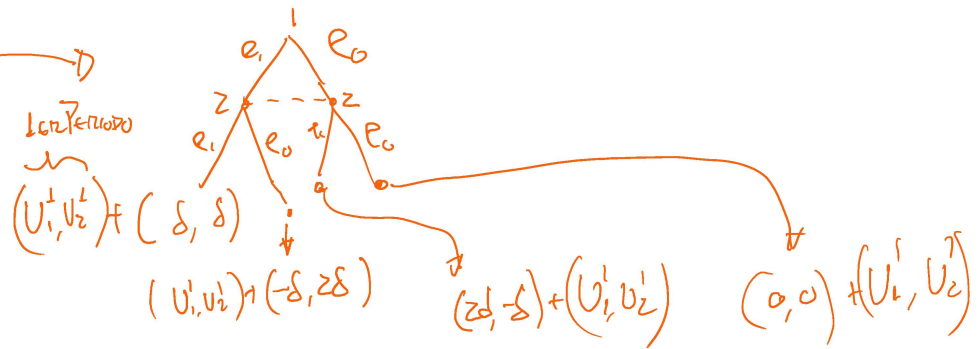
► Player 1 has 5 information sets in total

► A pure strategy for player 1 must specify what he does in each of these information sets



$$|S_1| = \left| \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ z & \delta & z & \delta & z \\ \delta & z & \delta & z & \delta \\ z & \delta & z & \delta & z \\ \delta & z & \delta & z & \delta \end{pmatrix} \right| = z^5 = 32$$

$$|S_2| = \left| \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ z & \delta & z & \delta & z \\ \delta & z & \delta & z & \delta \\ z & \delta & z & \delta & z \\ \delta & z & \delta & z & \delta \end{pmatrix} \right| = z^5 = 32$$



	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	$(\delta, \delta) + (U_1^t, U_2^t)$	$(-\delta, z\delta) + (U_1^t, U_2^t)$
$e_1 = 0$	$(z\delta, -\delta) + (U_1^t, U_2^t)$	$(0, 0) + (U_1^t, U_2^t)$

$$\Rightarrow (U_1^{t+1}, U_2^{t+1}) +$$

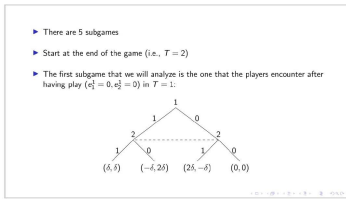
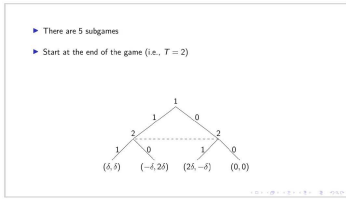
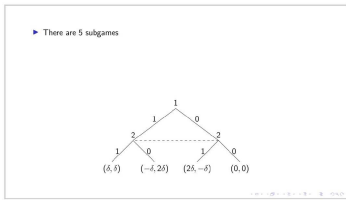
	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	δ, δ	$-\delta, z\delta$
$e_1 = 0$	$z\delta, -\delta$	$0, 0$

$$\Rightarrow (U_1, U_2)^{t+1} + \delta$$

	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	$1, 1$	$-1, z$
$e_1 = 0$	$z, -1$	$0, 0$

- We will solve for the set of pure SPNE of this game.
- Player 1 has 5 information sets in total
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- Player 1 has a total of $32 (2^5)$ pure strategies

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- Player 1 has 5 information sets in total
- A pure strategy for player 1 must specify what he does in each of these information sets
- Player 1 has a total of $32 (2^5)$ pure strategies
- Similarly, player 2 has a total of 32 pure strategies



The Nash equilibria can be seen by writing out the normal form of the game.

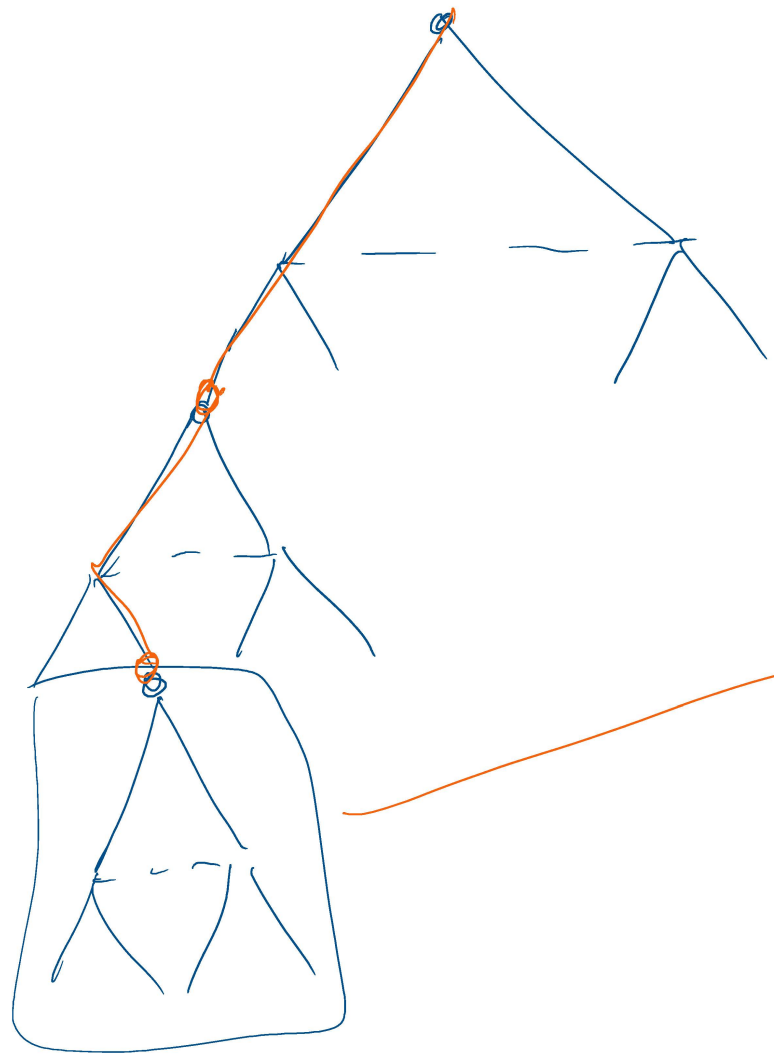
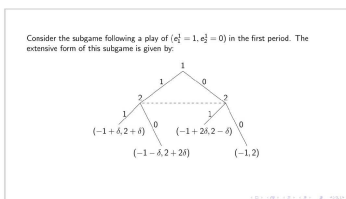
Normal Form of Extensive Form

$c_1 = 1$	$c_2 = 1$	$c_2 = 0$
$c_1 = 1$	δ, δ	$-1, 2\delta$
$c_1 = 0$	$2\delta, -\delta$	$0, 0$

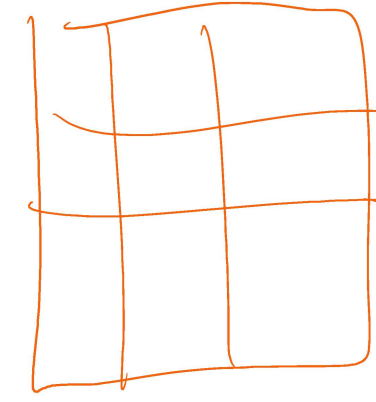
► This game has a unique Nash equilibrium in which the players play $(c_1^1 = 0, c_2^1 = 0)$

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► Therefore after having observed $(c_1^1 = 0, c_2^1 = 0)$ in the first period, both players will play $(c_1^2 = 0, c_2^2 = 0)$ in period 2



► $(U_1, U_2)^{t=1} + \delta(U_1, U_2)^{t=2} + \delta^2$



The normal form of this subgame can be seen in the Table

Normal Form of Extensive Form

	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	$-1, 1, 2 + \delta$	$-1, -1, 2 + 2\delta$
$e_1 = 0$	$-1 + 2\delta, 2 - \delta$	$-1, 2$

► $(e_1 = 0, e_2 = 0)$ is the unique Nash equilibrium

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► In any SPNE, $(e_1^2 = 0, e_2^2 = 0)$ must be played after observing $(e_1^1 = 1, e_2^1 = 0)$

► We can go through the remaining smaller subgames after the observation of $(e_1^1 = 1, e_2^1 = 0)$ and after the observation of $(e_1^1 = 1, e_2^1 = 1)$

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► We will reach the same conclusion in each of these scenarios: that $(e_1^2 = 0, e_2^2 = 0)$ must be played in each of these subgames

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► Why is this the case?

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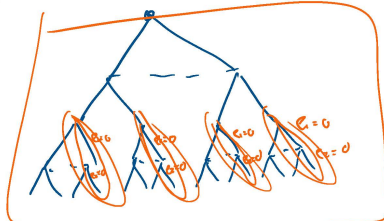
► Why is this the case?

► The idea is that payoffs that have accrued in period 1 are essentially sunk, and have no influence on incentives in period 2

To see this consider the normal form representation in the subgame after the observation of $(e_1^1 = 1, e_2^1 = 0)$

Normal Form of Extensive Form

	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	$-1, 1, 2 + \delta$	$-1, -1, 2 + 2\delta$
$e_1 = 0$	$-1 + 2\delta, 2 - \delta$	$-1, 2$



$T=1$

	$e_2 = 1$	$e_2 = 0$
$e_1 = 1$	$(1, 1) + U^{1-\delta} \delta$	$(-1, -1) + U^{1-\delta} \delta$
$e_1 = 0$	$(-1, 2) + U^{1-\delta} \delta$	$(0, 0) + U^{1-\delta} \delta$

\Rightarrow

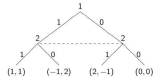
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- So what have we learned?
- Basically after any history, the strategic normal form is essentially the same as the original prisoner's dilemma
- Both players play $(a_1^t = 0, a_2^t = 0)$ after any information set in the last period

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- Now let us see what must be played in the first period by the two players
- Both players anticipate that $(a_1^2 = 0, a_2^2 = 0)$ will be played after any chosen action profile in the first period
- We can simplify the extensive form game to the following:



- If we draw the normal form of this game, then we get:

Normal Form of Extensive Form

	$a_2 = 1$	$a_2 = 0$
$a_1 = 1$	1,1	-1,2
$a_1 = 0$	2,-1	0,0

The unique Nash equilibrium of the above normal form game is $(a_1^1 = 0, a_2^1 = 0)$

- Therefore the unique SPNE is:
- $$\left(\begin{matrix} a_1^1 = 0 \\ a_2^1 = 0 \\ a_1^2 = 0 \\ a_2^2 = 0 \end{matrix} \right), \left(\begin{matrix} a_1^1 = 0 \\ a_2^1 = 0 \\ a_1^2 = 0 \\ a_2^2 = 0 \end{matrix} \right)$$
- In other words both players always shirk

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- ▶ Thus, the equilibrium outcome is simply the repetition of the unique NE of the stage game
- ▶ This holds more generally when the stage game has a unique NE
- ▶ Whenever the stage game has a unique NE, then the only SPNE of a finite horizon repeated game with that stage game is the repetition of the stage game NE

Theorem
 Suppose that the stage game G has exactly one NE, (s_1^*, \dots, s_n^*) . Then for any $\delta \in (0, 1)$ and any T the T -period repeated game has a unique SPNE in which all players play s_i^* at all information sets.
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- ▶ Knowing that the stage game Nash equilibrium is going to be played tomorrow, at any information set, we can ignore the past payoffs
- ▶ We concentrate just on the payoffs in the future. Thus in period $T-1$, player i simply wants to maximize:

$$\max_{A_i} \delta^{T-2} u_i(A_i, s_{-i}^{T-1}) + \delta^{T-1} u_i(A_i^*)$$

- ▶ What player i plays today has no consequences for what happens in period T since we saw that all players will play a^* no matter what happens in period $T-1$

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- ▶ So, the maximization problem above is the same as:

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- ▶ Following exactly this induction, we can conclude that every player must play a_i^* at all times and all histories

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