

Lecture 17: Applications of Subgame Perfect Nash Equilibrium

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- A game can repeat itself several times
- Static games turn into dynamic by repetition
- We will use  $(G, T)$  to denote that game  $G$  is repeated  $T$  times

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- Players observe the actions chosen for the players in period 1. Then in period 2, players simultaneously play the game  $G$ .
- This game proceeds until time  $T$ .
- After time  $T$ , if the action profiles chosen in times  $1, 2, \dots, T$  are given by  $(a_1^1, a_2^1), \dots, (a_1^T, a_2^T)$  then the total payoff is given by  $\sum_{t=1}^T \delta^{t-1} u_i(a_1^t, a_2^t)$

Consider the following two-player game:

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- This:  $u_i(a_1, a_2) = 2a_{-i} - a_i$

Prisoner's Dilemma (Game G)

	$a_2 = 1$	$a_2 = 0$
$a_1 = 1$	1, 1	0, 2
$a_1 = 0$	2, 0	0, 0

Best (1, 0) (0, 0)

- What happens when  $T = 1$

- What happens when  $T = 1$
- NE: Players 1 and 2 will both shirk ( $a_1 = 0, a_2 = 0$ )

Imagine players are engaged in a long run relationship that lasts more than just playing the game once:  $(G, T)$

- Both players play the simultaneous move game  $G$ .

$$\delta = \frac{1}{1+r}$$

$0 = \frac{1}{1+r}$

Imagine players are engaged in a long run relationship that lasts more than just playing the game once. (2,2)

1. Both players play the simultaneous move game G.

2. Both players observe the actions chosen by the two players. Then they play G again.

3. Then payoffs are calculated as the discounted sum of the utilities of the actions in each period with discount factor  $\delta \in (0,1)$ .

Suppose that the two players chose  $(a_1 = 1, a_2 = 1)$  in the first period in the second period, they chose  $(a_1 = 0, a_2 = 1)$ .

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Similarly, player 2 has a total of 32 pure strategies.

There are 5 subgames.

Start at the end of the game ( $t = T-2$ ).

The first subgame that we will analyze is the one that the players encounter after having played  $(a_1^t = 0, a_2^t = 0)$  in  $T-3$ .

The Nash equilibria can be seen by writing out the normal form of the game.

This game has a unique Nash equilibrium in which the players play  $(a_1^t = 0, a_2^t = 0)$ .

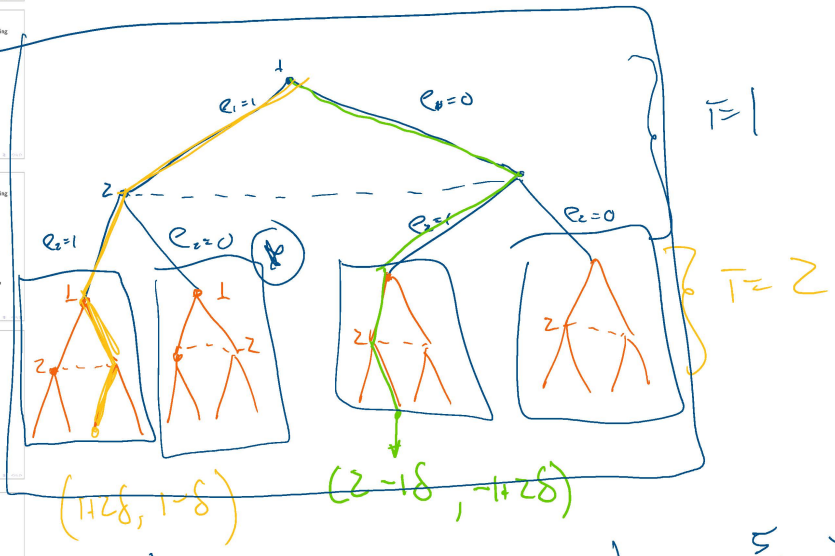
This game has a unique Nash equilibrium in which the players play  $(a_1^t = 0, a_2^t = 0)$ .

Therefore after having observed  $(a_1^t = 0, a_2^t = 0)$  in the first period, both players will play  $(a_1^{t+1} = 0, a_2^{t+1} = 0)$  in period  $t+1$ .

Consider the subgame following a play of  $(a_1^t = 1, a_2^t = 0)$  in the first period. The normal form of this subgame is given by:

The normal form of this subgame can be seen in the Table.

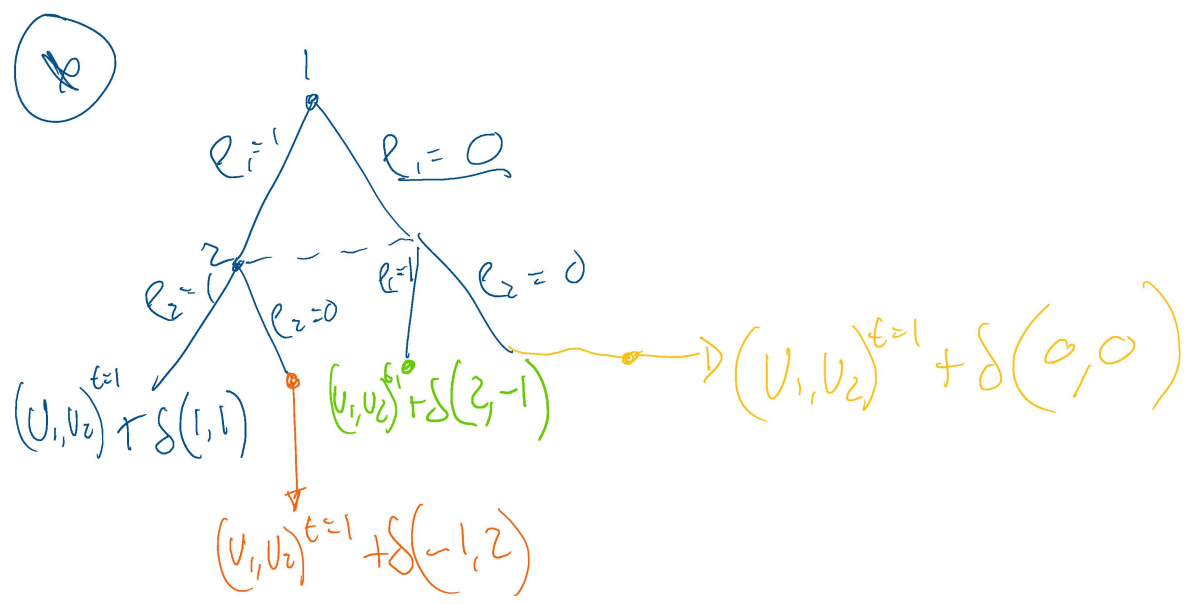
$(a_1 = 0, a_2 = 0)$  is the unique Nash equilibrium.



$(1,1, \delta, 1-\delta)$        $(2, -1, \delta, -1+\delta)$

$|S_1| = | \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 2 & 2 & 2 \end{pmatrix} | = 2^5 = 32$

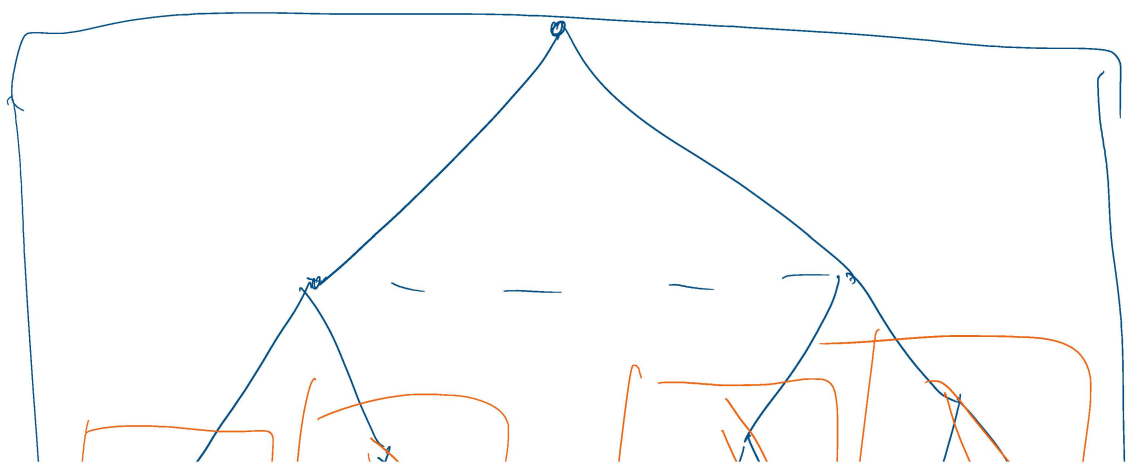
$|S_2| = 2^5 = 32$



	$R_2 = 1$	$R_2 = 0$
$R_1 = 1$	$(U_1, U_2)^{t+1} + \delta(1,1)$	$(U_1, U_2)^{t+1} + \delta(-1,2)$
$R_1 = 0$	$(U_1, U_2)^{t+1} + \delta(2,-1)$	$(U_1, U_2)^{t+1} + \delta(0,0)$

$(U_1, U_2)^{t+1} + \delta \begin{pmatrix} R_2 = 1 & R_2 = 0 \\ R_1 = 1 & 1, 1 & -1, 2 \\ R_1 = 0 & 2, -1 & 0, 0 \end{pmatrix}$

$\Rightarrow$  EN. SUBSUEGOS DE  $T=2$  ES  $(R_1=0, R_2=0)$



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$(a_1 = 0, a_2 = 0)$  is the unique Nash equilibrium.

- $(x_1 = 0, x_2 = 0)$  is the unique Nash equilibrium

► In any SPNE,  $(x_1^* = 0, x_2^* = 0)$  must be played after observing  $(x_1 = 1, x_2 = 0)$
- We can go through the remaining smaller subgames after the observation of  $(x_1 = 1, x_2 = 0)$  and after the observation of  $(x_1 = 1, x_2 = 1)$
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► We will reach the same conclusion in each of these scenarios: that  $(x_1^* = 0, x_2^* = 0)$  must be played in each of these subgames
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► Why is this the case?
- We can go through the remaining smaller subgames after the observation of  $(x_1 = 1, x_2 = 0)$  and after the observation of  $(x_1 = 1, x_2 = 1)$

► We will reach the same conclusion in each of these scenarios: that  $(x_1^* = 0, x_2^* = 0)$  must be played in each of these subgames

► Regardless of the observed action,  $(0, 0)$  is played in period 2

► Why is this the case?

► The idea is that payoffs that have accrued in period 1 are essentially sunk, and have no influence on reactions in period 2
- To see this consider the normal form representation in the subgame after the observation of  $(x_1 = 1, x_2 = 0)$

Normal Form of Extension Form

	$x_2 = 1$	$x_2 = 0$
$x_1 = 1$	$1, 1$	$-1, 2$
$x_1 = 0$	$2, -1$	$0, 0$
- We can subtract off the payoff that player 1 received in period 1 and divide through player 1's payoffs by 2 to obtain the following payoff matrix

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► We can do the same thing for player 2's payoffs and get the payoff matrix

Normal Form of Extension Form

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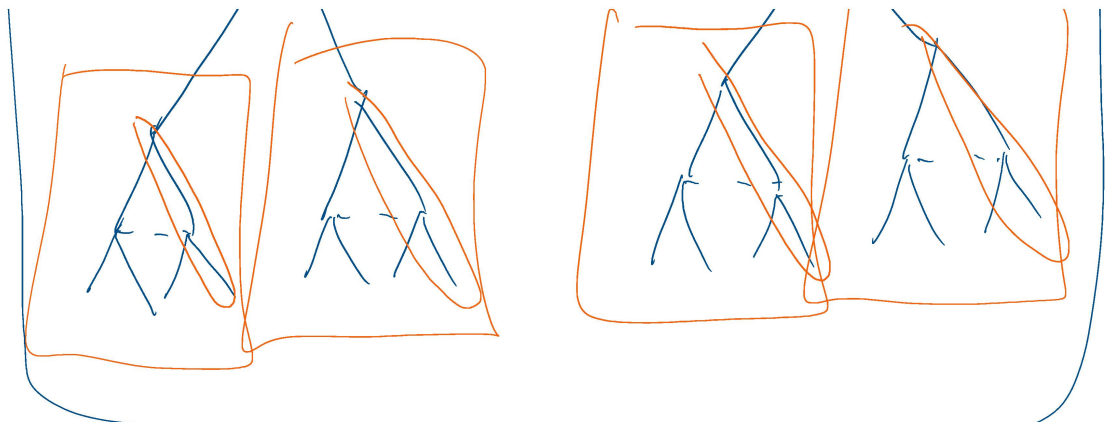
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► This normal form is just the original prisoner's dilemma

► This will be true no matter the action profile played in period 1
- So what have we learned?
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► Basically after any history, the strategic normal form is essentially the same as the original prisoner's dilemma

► Both players play  $(x_1^* = 0, x_2^* = 0)$  after any information set in the last period
- Now let us see what must be played in the first period by the two players



$$\bar{T} = 1$$

	$x_2 = 1$	$x_2 = 0$
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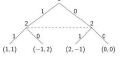
$$+ \delta(0, 0)$$

$$EN \Rightarrow (x_1 = 0, x_2 = 0)$$

- Now let us see what must be played in the first period by the two players
- Both players anticipate that  $(c^1 = 0, c^2 = 0)$  will be played after any chosen action profile in the first period

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- Now let us see what must be played in the first period by the two players
- Both players anticipate that  $(c^1 = 0, c^2 = 0)$  will be played after any chosen action profile in the first period
- We can simplify the extensive form game to the following:



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If we draw the normal form of this game, then we get:

	$a_1 = 1$	$a_1 = 0$
$a_2 = 1$	1, 1	2, 0
$a_2 = 0$	2, -1	0, 0

The unique Nash equilibrium of the above normal form game is  $(c^1 = 0, c^2 = 0)$

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Therefore the unique SPNE is:

$$\left( \begin{matrix} a_1^1 = 0 \\ a_1^2 = 0 \\ a_2^1 = 0 \\ a_2^2 = 0 \end{matrix} \right) \left( \begin{matrix} a_1^1 = 0 \\ a_1^2 = 0 \\ a_2^1 = 0 \\ a_2^2 = 0 \end{matrix} \right)$$

In other words both players always stick

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- Thus, the equilibrium outcome is simply the repetition of the unique NE of the stage game
- This holds more generally when the stage game has a unique NE
- Whenever the stage game has a unique NE, then the only SPNE of a finite horizon repeated game with that stage game is the repetition of the stage game NE

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**Theorem:**  
Suppose that the stage game  $G$  has exactly one NE  $(c_1^*, c_2^*, \dots, c_n^*)$ . Then for any  $T \in \mathbb{N}$  and any  $\Gamma$ , the  $T$ -times repeated game has a unique SPNE in which all players  $i$  play  $c_i^*$  at all information sets.

*new words los periodos!*

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- Knowing that the stage game Nash equilibrium is going to be played tomorrow, at any information set, we can ignore the past payoffs
- We concentrate just on the growth in the future. Thus in period  $T-1$ , player  $i$  simply wants to maximize:
 
$$\max_{a_i^{T-1}} \delta^T v_i(a_i^{T-1}, a_{-i}^{T-1}) + \delta^{T-1} v_i(a_i^{T-1})$$

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- What player  $i$  plays today has no consequences for what happens in period  $T$  since we saw that all players will play  $a_i^T$  no matter what happens in period  $T-1$

So, the maximization problem above is the same as:

$$\max_{a_i^{T-1}} v_i(a_i^{T-1}, a_{-i}^{T-1})$$

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Thus again, for this to be a Nash equilibrium, we need  $a_i^{T-1} = a_i^T, \dots, a_i^{T-1} = a_i^T$

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- What player  $i$  plays today has no consequences for what happens in period  $T$  since we saw that all players will play  $c^i$  no matter what happens in period  $T-1$ .
- So, the maximization problem above is the same as:
$$\max_{c^i} u_i(c^i, c_{-i}^{T-1})$$
- Thus again, for this to be a Nash equilibrium, we need  $c^1 = c_1^1, \dots, c^N = c_1^N$ .
- Following exactly this induction, we can conclude that every player must play  $c^i$  at all times and all histories.